

## (Improved) Optimal Triangulation of Saddle Surfaces

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## Motivation

- Smooth surface is locally approximated by a quadratic patch.
- Euclidean motion transforms the quadratic patch to graph of a bi-variate polynomial.
- $\rightarrow$ approximate graphs of quadratic polynomials!

$$
\{(x, y, z): z=F(x, y)\}
$$



- H. Pottmann, R. Krasauskas, B. Hamann, K. Joy, and W. Seibold: On piecewise linear approximation of quadratic functions. Journal for Geometry and Graphics 4 (2000), 31-53.

Introduction

## Interpolating Approximation

Non-interpolating Approximation

## Vertical Distance

- We are interested in a neighborhood of some point.
- Make the surface normal vertical.
- The direction in which Hausdorff distance is measured becomes almost vertical.


## Definition (Vertical Distance, $L_{\infty}$ Distance)

Given two domains $D_{1}, D_{2} \subset \mathbb{R}^{2}$ and two graphs $f: D_{1} \rightarrow \mathbb{R}$ and $g: D_{2} \rightarrow \mathbb{R}$ then the vertical distance is

$$
\operatorname{dist}_{V}(f, g)=\max _{(x, y) \in D_{1} \cap D_{2}}|f(x, y)-g(x, y)|
$$

## Properties of V-Distance

## Lemma

Let $A, B \subset \mathbb{R}^{3}$ be two sets with equal projection to the plane. Then

$\operatorname{dist}_{H}(A, B) \leq \operatorname{dist}_{V}(A, B)$



## V-Distance of Quadratic Functions

## Lemma (Every two points are the same)

Let $S$ be the graph of a quadratic function.
For every point $p \in S$, there is an affine transformation $\mathcal{T}_{p}: \mathbb{R}^{3} \rightarrow \mathbb{R}^{3}$ which satisfies the following:

- $\mathcal{T}_{p}(p)=\overrightarrow{0}$
- $\mathcal{T}_{p}(S)=$ a quadratic graph $\tilde{S}$ with a homogeneous polynomial of the form

$$
\begin{equation*}
\tilde{F}(x, y)=a x^{2}+b x y+c y^{2} \tag{*}
\end{equation*}
$$

- For all $q, r \in \mathbb{R}^{3}$ on a vertical line,

$$
|q-r|=\left|\mathcal{T}_{p}(q)-\mathcal{T}_{p}(r)\right| .
$$

- $\mathcal{T}_{p}(p)$ on the first two coordinates is a translation in $\mathbb{R}^{2}$.


## Vertical Distance of a Chord

If $S$ is negatively curved, the maximum distance to a triangle never occurs in the interior.

## Lemma

For a line segment $\overline{p q}$ between two points $p=\left(p_{x}, p_{y}, p_{z}\right)$ and $q=\left(q_{x}, q_{y}, q_{z}\right)$ on a quadratic graph $S$,

$$
\operatorname{dist}_{v}(\overline{p q}, S)=\frac{1}{4}\left|\tilde{F}\left(q_{x}-p_{x}, q_{y}-p_{y}\right)\right|
$$

- $\tilde{F}(x, y)$ is the homogeneous polynomial (*).
- The max. vertical distance is attained at the midpoint.



## Setup

From now on,

$$
S=\{(x, y, z): z=x y\}
$$

(by a linear transformation of the $x-y$-plane)

## Goal

Given $\varepsilon>0$, find a triangle $T$ with vertices $p_{0}, p_{1}, p_{2} \in S$ of largest area such that

$$
\operatorname{dist}_{V}(T, S) \leq \varepsilon
$$

Translated and reflected copies of $T$ have the same error and tile the plane:
max. AREA $\Leftrightarrow$ min. NUMBER of triangles

## Maximize the Area of Planar Triangles



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## Triangulate the Saddle

Lift the planar triangulation to the surface


## Can We Do Better?

## What do we have?

Given an $\varepsilon>0$ and a saddle surface $S$, we can find a family $\mathcal{T}$ of triangles which interpolate the surface and

- have maximum area,
- maintain $\operatorname{dist}_{V}(S, T) \leq \varepsilon$ for all $T \in \mathcal{T}$.


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## Question. . .

- Can this be improved by allowing non-interpolating triangles?
- Pottmann et al. (2000) conjectured NO.

This question is easy for convex approximation.

## Pseudo-Euclidean Transformations

- A $\lambda$-pseudo Euclidean map is given by:

$$
(x, y) \mapsto\left(\lambda x, \frac{1}{\lambda} y\right)
$$

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- Area (projected) is preserved.
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- one-parameter family of area preserving triangles
- How should they be lifted?



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- Lift the triangle vertically such that the distance to $S$ is minimized.



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- Lift vertices off the surface by $\alpha$ :

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S_{\alpha}=\{(x, y, z): z=x y+\alpha\}
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- Lift the triangle vertically such that the distance to $S$ is minimized.
- Lift vertices off the surface by $\alpha$ :

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- Vertical distance is attained at midpoints.



## Vertical Perturbed Lifting (Cont.)

- Vertical distances from edges to $S$ are

$$
\begin{aligned}
& \frac{\xi \eta}{4}+\alpha>0 \\
& \frac{1}{4}(\xi-\eta)^{2}-\alpha>0
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and have to be equal.


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- $\alpha=\frac{1}{8}\left(\xi^{2}-3 \xi \eta+\eta^{2}\right)$



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- and the vertical distance is

$$
\frac{\sqrt{15}}{4} \varepsilon \approx 0.968246 \varepsilon
$$



## Picture in Space



## The Planar Super-Optimal Triangle

- Pseudo-euclidean motions give a one-parameter family of optimal triangles.



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- Note the tangency property



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## OPEN: <br> Lift vertices by different amounts?



