

Testing Congruence of Point Sets

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Testing Congruence of Point Sets and Symmetry Günter Rote Freie Universität Berlin



Overview



- 1 dimension
- 2 dimensions
- 3 dimensions
- 4 dimensions
- *d* dimensions

Overview



- 1 dimension
- 2 dimensions
- 3 dimensions
- 4 dimensions ?
- d dimensions $O(n^{\lceil d/3 \rceil} \log n)$ time

 $O(n \log n)$ time

Overview



- 1 dimension
- 2 dimensions
- 3 dimensions
- 4 dimensions ?
- d dimensions $O(n^{\lceil d/3 \rceil} \log n)$ time
- Problem statement and variations
- PRUNING and DIMENSION REDUCTION
- Point groups (discrete subgroups of the orthogonal group)

 $O(n \log n)$ time

Rotation or Rotation+Reflection?

We only need to consider *proper* congruence (orientation-preserving congruence, of determinant +1).

If mirror-congruence is also desired, repeat the test twice, for B and its mirror image B'.



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Congruence = Rotation + Translation

Translation is easy to determine:

The centroid of A must coincide with the centroid of B.



 \rightarrow from now on: All point sets are centered at the origin O:

$$\sum_{a \in A} a = \sum_{b \in B} b = 0$$

We need to find a rotation around the origin (orthogonal matrix T with determinant +1) which maps A to B.

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Exact Arithmetic

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The proper setting for this (mathematical) problem requires real numbers as inputs and exact arithmetic.

 \rightarrow the *Real RAM* model (RAM = random access machine): One elementary operation with real numbers (+, \div , $\sqrt{}$, sin) is counted as one step.



Applications



Congruence testing is the basic problem for many pattern matching tasks

- computer vision
- star matching
- brain matching
- . . .

The proper setting for this applied problem requires tolerances, partial matchings, and other extensions.

Approximate matching



Given two sets A and B in the plane and an error tolerance ε , find a bijection $f: A \to B$ and a congruence T such that



This problem is NP-hard. [S.Iwanowski 1991, C.Dieckmann 2012]

Arbitrary Dimension



$$A, B \subset \mathbb{R}^d$$
, $|A| = |B| = n$.

We consider the problem for fixed dimension d.

When d is unrestricted, the problem is equivalent to graph isomorphism:

$$G = (V, E), V = \{1, 2, \dots, n\}$$

$$\mapsto A = \underbrace{\{e_1, \dots, e_n\}}_{\text{regular simplex}} \cup \{\frac{e_i + e_j}{2} \mid ij \in E\} \subset \mathbb{R}^n$$

CONJECTURE:

Congruence can be tested in $O(n \log n)$ time for every fixed dimension d. ("fixed-parameter tractable") Current best bound: $O(n^{\lceil d/3 \rceil} \log n)$ time

One dimension

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Trivial.

(after shifting the centroid to the origin and getting rid of reflection):

Test if A = B. $O(n \log n)$ time.

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Testing Congruence and Symmetry of Point Sets

Two dimensions

Can be done by string matching. Sort points around the origin.

Encode alternating sequence of distances r_i and angles ϕ_i .



 $(r_1, \phi_1, r_2, \phi_2, \dots, r_n, \phi_n)$





[Manacher 1976]

Three dimensions

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Compute the convex hull P(A) and P(B), in $O(n \log n)$ time.

Check isomorphism between the corresponding planar graphs, in O(n) time. [Hopcroft and Wong 1974]

The result is unique, up to

- the symmetries of a Platonic solid (at most 60 choices), or
- a rotation around an axis.

 \rightarrow project down to a 2-dimensional problem.

[Sugihara 1984; Alt, Mehlhorn, Wagener, Welzl 1988]







Pruning





Find *some* criterion that distinguishes points (distance from the center, number of closest neighbors, ...)



Pruning





Simultaneously, apply the same pruning procedure to B.

Dimension Reduction



As soon as |A'| = |B'| = k is small: Choose a point $a_0 \in A'$ and try all k possibilities of mapping it to a point $b \in B'$.

Fixing $a_0 \mapsto b$ reduces the dimension by one.



Project perpendicular to Oa_0 and label projected points a'_i with the signed projection distance d_i as (a'_i, d_i) .



PRUNING:

Find some (geometric, combinatorial) characteristic that distinguishes points from each other. Keep only the smallest equivalence class.

DIMENSION REDUCTION:

Reduce one *d*-dimensional problem to k problems of dimension d-1.

[M. D. Atkinson, J. Algorithms 1987, for d = 3]

Three Dimensions

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If the points lie in a plane or on a line \rightarrow DIMENSION REDUCTION.

Compute the convex hull P(A). If there are vertices of different degrees \rightarrow PRUNE

The number n of vertices is reduced to $\leq n/2$. RESTART. All n vertices have now degree 3, 4, or 5. There are $f = \frac{n}{2} + 2$ or f = n + 2 or $f = \frac{3n}{2} + 2$ faces.

If the face degrees are not all equal \rightarrow switch to the centroids of the faces and PRUNE them. n is reduced to $\leq \frac{3n}{4} + 1$. RESTART.

Now P(A) must have the graph of a Platonic solid. $\rightarrow n \leq 20$. \rightarrow DIMENSION REDUCTION.

Three Dimensions

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If the points lie in a plane or on a line \rightarrow DIMENSION REDUCTION.

Compute the convex hull P(A). $\checkmark O(|A| \log |A|)$ time If there are vertices of different degrees \rightarrow PRUNE

The number n of vertices is reduced to $\leq n/2$. RESTART. All n vertices have now degree 3, 4, or 5. There are $f = \frac{n}{2} + 2$ or f = n + 2 or $f = \frac{3n}{2} + 2$ faces.

If the face degrees are not all equal \rightarrow switch to the centroids of the faces and PRUNE them. n is reduced to $\leq \frac{3n}{4} + 1$. RESTART.

$$TIME = O(n \log n) + O(\frac{3}{4}n \log \frac{3}{4}n) + O((\frac{3}{4})^2 n \log((\frac{3}{4})^2 n)) + \dots = O(n \log n)$$

Symmetry groups



COROLLARY. The symmetry group of a finite full-dimensional point set in 3-space (= a discrete subgroup of O(3)) is

- the symmetry group of a Platonic solid,
- the symmetry group of a regular prism,
- or a subgroup of such a group.



The *point groups* (discrete subgroups of O(3)) are classified (Hessel's Theorem). [F. Hessel 1830, M. L. Frankenheim 1826]

Point groups in higher dimensions



Is this true in higher dimensions?

¿The symmetry group of a finite full-dimensional point set in d-space (= a discrete subgroup of O(d)) is

- the symmetry group of a regular *d*-dimensional polytope:
 - a regular simplex
 - * a hypercube (or its dual, the crosspolytope)
 - a regular n-gon in two dimensions
 - a dodecahedron (or its dual, the icosahedron) in 3 d.
 - a 24-cell, or a 120-cell (or its dual, the 600-cell) in 4 d.
- the symmetry group of the Cartesian product of lower-dimensional regular polytopes,
- or a subgroup of such a group?

The symmetry groups of the root systems E_6 , E_7 , E_8 in 6, 7, and 8 dimensions might be counterexamples.

Dimension d

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Dimension reduction without pruning:

Pick $a_0 \in A$. Try $a_0 \mapsto b$ for all $b \in B$ (*n* possibilities). $\rightarrow O(n^{d-2} \log n)$ time

Improvement [Matoušek \approx 1998]: $\overset{a_0}{\bullet} \overset{a_1}{\bullet}$ Consider all *closest pairs* of *A* and *B*. Each point belongs to

 $\leq C_d$ closest pairs. (packing argument, the kissing number). $\implies O(n)$ closest pairs.

Pick a closest pair $a_0a_1 \in A$. Try $(a_0, a_1) \mapsto (b, b')$ for all closest pairs $(b, b') \in B$. O(n) possibilities, reducing the dimension by *two*. $\rightarrow O(n^{\lfloor d/2 \rfloor} \log n)$ time

Further improvement: Find a "closest triplet" ...





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Compute the closest pair graph

$$G(A) = (A, \{ aa' : ||a - a'|| = \delta \})$$

where δ is the distance of the closest pair.

By the PRUNING principle, we can assume that all points look locally the same:

• same distance from the origin. (A lies on the 3-sphere \mathbb{S}^3 .)

 All points have congruent neighborhoods in G(A). (The neighbors of a lie on a 2-sphere in S³; There are at most 12 neighbors.)

CASE 1. The vertex figure has no symmetries.



CASE 1.

The vertex figure has no symmetries.

Follow the sequence of "1"-neighbors



 a_1

*a*₀ 1____





CASE 1.

 a_0

The vertex figure has no symmetries.

Follow the sequence of "1"-neighbors

 a_1

• If the "1"-edges form a matching: Take the midpoints of these edges. $n \rightarrow n/2$.





CASE 1.

The vertex figure has no symmetries.

Follow the sequence of "1"-neighbors

 a_1

$(a_0, a_1, a_2) \cong (a_1, a_2, a_3) \cong (a_2, a_3, a_4) \cong \cdots$ A is partitioned into cycles of the same length.

- Cycles are short \rightarrow their centroid is nonzero; replace them by their centroids.
- a₀, a₁, a₂ lie on a circle through the origin special situation: all these circles are parallel; or there is a bounded number of circles.
- a_0, a_1, a_2 span a hyperplane \rightarrow take their normals





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CASE 2.

The vertex figure has, say, tetrahedral symmetry.

```
Start with a path of length 3,
by extending a_0a_1 "as straight as possible"
(Finitely many starting patterns).
a_0, a_1, a_2 span a hyperplane.
```



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CASE 2.

The vertex figure has, say, tetrahedral symmetry.

```
Start with a path of length 3,
by extending a_0a_1 "as straight as possible"
(Finitely many starting patterns).
a_0, a_1, a_2 span a hyperplane.
```



Extend the path by

$$(a_0, a_1, a_2) \cong (a_1, a_2, a_3) \cong (a_2, a_3, a_4) \cong \cdots$$

 $\rightarrow A$ is partitioned into cycles of the same length.

- [W. Threlfall and H. Seifert, Math. Annalen, 1931, 1933] enumerated discrete subgroups of SO(4) (determinant +1)
- [J. Conway and D. Smith 2003] complete enumeration of point groups
- 4d-rotation $T \leftrightarrow \text{pair}(R, S)$ of 3d-rotations. (for example, via quaternions)
- Goursat's Lemma: [É. Goursat 1890] Pairs of 3d point groups + additional information
- \rightarrow 4d point groups



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 The groups generated by reflections (Coxeter groups) have been enumerated up to 8 dimensions.
 [Norman Johnson, unpublished book manuscript]

umerate the groups		Generators (See	section 3	able 4.1.	The chira	al groups
Grou	$\frac{ip}{O}$ [<i>i</i> _I ,	$1], [\omega, 1], [1, i_O], [1, \omega]$]; (groups of		0 1
$\pm [I \times$	O $[i_I, [i_I]]$	1], $[\omega, 1], [1, i], [1, \omega]$;			
$\pm [I \times$	T $[i_T]$	1] $[\omega, 1], [1, e_n], [1, j]$]; O	rientation	n-preservir	າg
$\pm [I \times$	D_{2n}] $[\iota_I,$	$i_1, [\omega, -], [\nu, 1], [1, e_n];$	0	rthogonal	l transform	mations)
$\pm [I \times$	C_n]	$[1, 1], [\omega, 1], [-j, -n],$:	rtiogona		nationsj
$\pm [O \times$	T] [i_O	$, 1], [\omega, 1], [1, o], [1, \omega]$:]•	r a		
$\pm [O \times$	$[D_{2n}]$ $[i_O$	$, 1], [\omega, 1], [1, e_n], [1, J]$	(] 9 ;]	[Conwa	y and Smi	th 2003 J
$\pm \frac{1}{2}[O \times$	$[D_{2n}]$ [<i>i</i> ,	$1], [\omega, 1], [1, e_n]; [i_0, j]$	/] ₁	_		_
$\pm \frac{1}{2}[O \times$	$(\overline{D}_{4n}] \qquad [i,1],[a]$	$[\omega, 1], [1, e_n], [1, j]; [i_0]$	$,e_{2n}]$			
$\pm \frac{1}{6} [O >$	$< D_{6n}$] [i, 1], [$j,1],[1,e_{n}];[i_{O},j],[\omega$	$[,e_{3n}]$			
$\pm [O >$	$\langle C_n]$	$[i_O, 1], [\omega, 1], [1, e_n];$				
$\pm \frac{1}{2}[O > $	$\times C_{2n}$] [<i>i</i> , 1	$[], [\omega, 1], [1, e_n]; [i_O, e_n]$	$_{2n}]$			
$\pm[T:$	$\times D_{2n}$] [i,	$[1], [\omega, 1], [1, e_n], [1, j]$];			
$\pm [T]$	$\times C_n$]	$[i,1], [\omega,1], [1,e_n];$				
$\pm \frac{1}{3}[T$	$\times C_{3n}$]	$[i,1], [1,e_n]; [\omega,e_{3n}]$				
$\pm \frac{1}{2}[D_{2m}$	$\times \overline{D}_{4n}$] [e_n	$[1, 1], [1, e_n], [1, j]; [j, e_n]$	2n			
$\pm [D_{2m}$	$\times C_n$]	$[e_m, 1], [j, 1], [1, e_n];$				
$\pm \frac{1}{2}[D_{2m}$	$\times C_{2n}$]	$[e_m, 1], [1, e_n]; [j, e_{2n}]$	nie en			
$+\frac{1}{2}[D_{2m}$	$\times C_{2n}$]	· +	lor	- both m	and n must	· he odd
$\pm \frac{1}{2}[\overline{D}_{4m}]$	$(\times C_{2n}]$ [e _m	$,1],[j,1],[1,e_n];[e_{2m},$	$e_{2n}]$			
Table 4.1. Chiral gro	ups 1 Those					
4.6—some others app	bear in the last few	ost of the "metachiral" g	groups—s y of Po	int Sets	,CSA2D	Koper, June 9–13, 2013
		able 4.2.				

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	Generators	Coxeter Name
Group	$[i, 1] [w, 1], [1, i_I], [1, w];$	$[3, 3, 5]^+$
$\pm [I \times I]$	$[i_{I}, 1], [w, 1], [-i_{I}, i_{I}]$	$2.[3,5]^+$
$\pm \frac{1}{60}[I \times I]$	(w, w), (v, v)	$[3, 5]^+$
$+\frac{1}{60}[I \times I]$	$i \downarrow i$	$2.[3, 3, 3]^+$
$\pm \frac{1}{60}[I \times \overline{I}]$	$;[\omega,\omega],[v_1,v_1]$	$[3, 3, 3]^+$
$+\frac{1}{60}[I \times \overline{I}]$; + , +	$[3, 4, 3]^+: 2$
$\pm [O \times O]$	$[i_{O}, 1], [\omega, 1], [1, i_{O}], [1, \omega],$	$[3, 4, 3]^+$
$\pm \frac{1}{2}[O \times O]$	$[i, 1], [\omega, 1], [1, i], [1, \omega]; [io, io]$	$[3, 3, 4]^+$
$\pm \frac{1}{6}[O \times O]$	$[i, 1], [j, 1], [1, i], [1, j]; [\omega, \omega], [i_0, i_0]$	[0, 0, 4]
$\pm \frac{1}{24}[O \times O]$	$;[\omega,\omega],[i_O,i_O]$	$2 \cdot [0, 4]$
$+\frac{1}{24}[O \times O]$; + , +	[0, 4]
$+\frac{1}{24}[O \times \overline{O}]$; + , -	[2, 3, 3]
$\pm [T \times T]$	$[i,1],[\omega,1],[1,i],[1,\omega];$	[3, 4, 3]
$\pm \frac{1}{3}[T \times T]$	$[i,1],[j,1],[1,i],[1,j];[\omega,\omega]$	$[^+3, 3, 4^+]$
$\cong \pm \frac{1}{3} [T \times \overline{T}]$	$[i,1],[j,1],[1,i],[1,j];[\omega,\overline{\omega}]$	"
$\pm \frac{1}{12}[T \times T]$	$;[\omega,\omega],[i,i]$	$2.[3,3]^+$
$\cong \pm \frac{1}{12} [T \times \overline{T}]$	$;[\omega,\overline{\omega}],[i,-i]$	"
$+\frac{1}{12}[T \times T]$; + , +	$[3,3]^+$
$\cong +\frac{1}{12}[T \times \overline{T}]$; + , +	"
$\pm [D_{2m} \times D_{2n}]$	$[e_m, 1], [j, 1], [1, e_n], [1, j];$	
$\pm \frac{1}{2} [\overline{D}_{4m} \times \overline{D}_{4n}]$	$[e_m,1],[j,1],[1,e_n],[1,j];[e_{2m},e_{2n}]$	
$\pm \frac{1}{4} [D_{4m} \times \overline{D}_{4n}]$	$[e_m,1],[1,e_n];[e_{2m},j],[j,e_{2n}]$	Conditions
$+\frac{1}{4}[D_{4m} \times D_{4n}]$	- , $-$; $+$, $+$	$m,n { m odd}$
$\pm \frac{1}{2f} [D_{2mf} \times D_{2nf}^{(s)}]$	$[e_m, 1], [1, e_n]; [e_{mf}, e_{nf}^s], [j, j]$	(s,f)=1
$+\frac{1}{2f}[D_{2mf} \times D_{2nf}^{(s)}]$	- , - ; + , +	m, n odd, (s, 2f) = 1
$\pm \frac{1}{f} \begin{bmatrix} C_{mf} \times C_{nf}^{(s)} \end{bmatrix}$	$[e_m, 1], [1, e_n]; [e_{mf}, e_{nf}^s]$	(s,f)=1
$+\overline{f}[C_{mf} \times C_{nf}]$; +	m, n odd, (s, 2f) = 1

Table 4.2. The *chiral* groups (continued)

Table 4.2. Chiral groups, II. These groups are mostly "orthochiral." with a few

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Group	Extending element	Coxeter Name
$\pm [I \times I] \cdot 2$	A THE R. P. LEWIS	[3, 3, 5]
$\pm \frac{1}{20}[I \times I] \cdot 2$	*	2.[3, 5]
$+\frac{1}{20}[I \times I] \cdot 2_3 \text{ or } 2_1$	* or*	[3,5] or [3,5]°
$\pm \frac{1}{60} [I \times \overline{I}] \cdot 2$	5	2.[3, 3, 3]
$+\frac{1}{50}[I \times \overline{I}] \cdot 2_3 \text{ or } 2_1$	* or - *	[3,3,3]° or [3,3,3]
$\pm [O \times O] \cdot 2$	*	[3, 4, 3]: 2
$\pm \frac{1}{2}[O \times O] \cdot 2 \text{ or } \overline{2}$	* or $* [1, i_0]$	$[3, 4, 3]$ or $[3, 4, 3]^+ 2$
$\pm \frac{1}{6}[O \times O] \cdot 2$	*	[3, 3, 4]
$\pm \frac{1}{24}[O \times O] \cdot 2$	* *	2.[3, 4]
$+\frac{1}{24}[O \times O] \cdot 2_3 \text{ or } 2_1$	* or*	$[3, 4]$ or $[3, 4]^{\circ}$
$+\frac{1}{24}[O \times \overline{O}] \cdot 2_3$ or 2_1	* or*	[2,3,3]° or [2,3,3]
$\pm [T \times T] \cdot 2$	*	$[3, 4, 3^+]$
$\pm \frac{1}{3}[T \times T] \cdot 2$	al a mi s* d a iran	[+3, 3, 4]
$\pm \frac{1}{3} [T \times \overline{T}] \cdot 2$	a de la seconda de	$[3, 3, 4^+]$
$\pm \frac{1}{12} [T \times T] \cdot 2$	*	2.[+3,4]
$\pm \frac{1}{12} [T \times \overline{T}] \cdot 2$	*	2.[3, 3]
$+\frac{1}{12}[T \times T] \cdot 2_3 \text{ or } 2_1$	* or - *	[⁺ 3, 4] or [⁺ 3, 4]°
$+\frac{1}{12}[T \times \overline{T}] \cdot 2_3 \text{ or } 2_1$	* or*	[3,3]° or [3,3]
$\pm [D_{2n} \times D_{2n}] \cdot 2$	*	and the set of the sectors
$\pm \frac{1}{2} [D_{4n} \times \overline{D}_{4n}] \cdot 2 \text{ or } \overline{2}$	* or $*[1, e_{2n}]$	
$\pm \frac{1}{4} [D_{4n} \times \overline{D}_{4n}] \cdot 2$	*	Conditions
$+\frac{1}{4}[D_{4n} \times D_{4n}] \cdot 2_3 \text{ or } 2_1$	* or - *	n odd
$\frac{1}{2f} \begin{bmatrix} D_{2nf} \times D_{2nf}^{(s)} \end{bmatrix} \cdot 2^{(\alpha,\beta)} \text{ or } \overline{2}$	$*[e_{2nf}^{\alpha}, e_{2nf}^{\alpha s + \beta f}] \text{ or } *[1, j]$	See
$\frac{2f[D_{2nf} \times D_{2nf}^{(s)}] \cdot 2^{(\alpha,\beta)} \text{ or } \overline{2}}{\pm \frac{1}{2}[C_{\alpha,\beta}] \cdot 2^{(\alpha,\beta)} \text{ or } \overline{2}}$	$*[e_{2nf}^{\alpha}, e_{2nf}^{\alpha s + \beta j}] \text{ or } *[1, j]$	Text
$= \int [C_{nf} \times C_{nf}^{(s)}] \cdot 2^{(\gamma)}$ $+ \frac{1}{2} [C_{nf} \times C_{nf}^{(s)}] \cdot 2^{(\gamma)}$	$*[1, e_{2nf}^{\gamma(f,s+1)}]$	in
$f(Onf \times Onf] \cdot 2^{(1)}$	$*[1, e_{2nf}^{\prime(j, (r+1))}]$) Appendix

Table 4.3. The *achiral* groups

Table 4.3. Achiral groups.

$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	S
$\pm \frac{1}{2} [I \times I] \cdot 2$ * 2.[3,5] I ne achiral group	S
$+\frac{1}{60}[I \times I] \cdot 2_3 \text{ or } 2_1 $ * or $-*$ [3,5] or [3,5]°	
$\pm \frac{1}{60}[I \times \overline{I}] \cdot 2$ * 2.[3, 3, 3]	
$+\frac{1}{60}[I \times \overline{I}] \cdot 2_3 \text{ or } 2_1 $ * or $-* [3,3,3]^\circ \text{ or } [3,3,3] $ • Visualize these	
$\pm [O \times O] \cdot 2$ * $[3, 4, 3] : 2$	
$\pm \frac{1}{2}[O \times O] \cdot 2 \text{ or } \overline{2}$ * or * [1, i _O] [3, 4, 3] or [3, 4, 3] ^{+·} 2 Groups:	
$\pm \frac{1}{6}[O \times O] \cdot 2$ * [3,3,4] Schlogel diagram	of o
$\pm \frac{1}{24}[O \times O] \cdot 2$ * 2.[3,4] Schleger diagram	JId
$+\frac{1}{24}[O \times O] \cdot 2_3 \text{ or } 2_1$ * or -* [3,4] or [3,4]° 4 -polytope which	has
$+\frac{1}{24}[O \times \overline{O}] \cdot 2_3 \text{ or } 2_1 $ * or $-*$ [2,3,3]° or [2,3,3] + polycope withen	nas
$\pm [T \times T] \cdot 2$ * $[3,4,3^+]$ these symmetries.	
$\pm \frac{1}{3}[T \times T] \cdot 2$ * [+3,3,4]	
$\pm \frac{1}{3} [T \times \overline{T}] \cdot 2 \qquad \qquad * \qquad [3, 3, 4^+]$	
$\pm \frac{1}{12}[T \times T] \cdot 2 $ * 2.[+3,4]	
$\pm \frac{1}{12}[T \times T] \cdot 2 $ * 2.[3,3]	
$+\frac{1}{12}[T \times T] \cdot 2_3 \text{ or } 2_1 $ * or $-*$ [+3,4] or [+3,4]°	
$+\frac{1}{12}[T \times T] \cdot 2_3 \text{ or } 2_1 $ * or $-*$ [3,3]° or [3,3]	
$\pm [D_{2n} \times D_{2n}] \cdot 2 \qquad \qquad *$	
$\begin{array}{c} \pm \frac{1}{2} [D_{4n} \times D_{4n}] \cdot 2 \text{ or } 2 \\ \pm \frac{1}{2} [D_{4n} \times \overline{D}_{4n}] \cdot 2 \text{ or } 2 \\ \end{array} \qquad \qquad$	
$+\frac{1}{4}[D_{4n} \times D_{4n}] \cdot 2 \qquad * \qquad \underbrace{Conditions}_{1}$	
$ \begin{array}{c} + 4[D_{4n} \times D_{4n}] \cdot 2_3 \text{ or } 2_1 \\ \pm \frac{1}{22}[D_{2n} \times D^{(s)}] + 2(\alpha, \beta) \\ \hline \end{array} \qquad \qquad$	
$\frac{2f(-2nf \wedge D_{2nf}) \cdot 2^{(-p)} \text{ or } 2}{\frac{1}{2} \left[D_{2nf} \times D^{(s)} \right] \cdot 2^{(\alpha,\beta)} - \overline{2}} \qquad \left[e_{2nf}^{\alpha,\beta} e_{2nf}^{\alpha,\beta} \right] \text{ or } \ast [1,j] \qquad \qquad \text{See}$	
$\pm \frac{1}{f} [C_{nf} \times C^{(s)}] \cdot 2^{(\gamma)} \qquad \qquad$	
$+\frac{1}{f} \begin{bmatrix} C_{nf} \times C_{nf}^{(s)} \end{bmatrix} \cdot 2^{(\gamma)} \qquad \qquad$	

Group	Extending element	Coxeter Name	Table 4.3.
$\pm [I \times I] \cdot 2$	1. A 1 (A 1 A 1 A 1 A 1 A 1 A 1 A 1 A 1 A	[3, 3, 5]	
$\pm \frac{1}{50}[I \times I] \cdot 2$	*	2.[3, 5]	The <i>achiral</i> groups
$+\frac{1}{60}[I \times I] \cdot 2_3 \text{ or } 2_1$	* or *	$[3, 5]$ or $[3, 5]^{\circ}$	
$\pm \frac{1}{60}[I \times \overline{I}] \cdot 2$	*	2.[3, 3, 3]	x /
$+\frac{1}{60}[I \times \overline{I}] \cdot 2_3 \text{ or } 2_1$	* or - *	[3, 3, 3]° or [3, 3, 3]	 Visualize these
$\pm [O \times O] \cdot 2$	*	[3, 4, 3]: 2	
$\pm \frac{1}{2}[O \times O] \cdot 2 \text{ or } \overline{2}$	* or $*[1, i_0]$	$[3, 4, 3]$ or $[3, 4, 3]^+ 2$	groups:
$\pm \frac{1}{6}[O \times O] \cdot 2$	*	[3, 3, 4]	Schlogol diagram of a
$\pm \frac{1}{24}[O \times O] \cdot 2$	*	2.[3, 4]	Schleger diagraffi Or a
$+\frac{1}{24}[O \times O] \cdot 2_3 \text{ or } 2_1$	* or*	$[3, 4]$ or $[3, 4]^{\circ}$	4-nolytone which has
$+\frac{1}{24}[O \times \overline{O}] \cdot 2_3 \text{ or } 2_1$	* or *	$[2, 3, 3]^{\circ}$ or $[2, 3, 3]$	+ polytope which has
$\pm [T \times T] \cdot 2$	*	$[3, 4, 3^+]$	these symmetries.
$\pm \frac{1}{3}[T \times T] \cdot 2$	and the second parts	[+3, 3, 4]	
$\pm \frac{1}{3}[T \times \overline{T}] \cdot 2$	*	$[3, 3, 4^+]$	Then go to 5d and
$\pm \frac{1}{12} [T \times T] \cdot 2$	*	2.[+3,4]	
$\pm \frac{1}{12} [T \times T] \cdot 2$	*	2.[3, 3]	higher
$+\frac{1}{12}[T \times T] \cdot 2_3 \text{ or } 2_1$	* or - *	[⁺ 3,4] or [⁺ 3,4]°	ingliei
$+\frac{1}{12}[T \times T] \cdot 2_3 \text{ or } 2_1$	* or*	$[3,3]^{\circ}$ or $[3,3]$	
$\pm [D_{2n} \times D_{2n}] \cdot 2$	*		
$\pm \frac{1}{2} [D_{4n} \times D_{4n}] \cdot 2 \text{ or } 2$ $\pm \frac{1}{2} [D_{4n} \times \overline{D}_{4n}] \cdot 2 $	* or $* [1, e_{2n}]$		
$+\frac{1}{4}[D_{4n} \times D_{4n}] \cdot 2$ $+\frac{1}{2}[D_{4n} \times \overline{D}_{4n}] \cdot 2$	*	Conditions	
$\pm \frac{1}{24} [D_{2n} \times D_{4n}] \cdot 2_3 \text{ or } 2_1$ $\pm \frac{1}{24} [D_{2n} \times D^{(s)}] + 2(\alpha, \beta) = \overline{\alpha}$	* or - *	n odd	
$\frac{2f(-2nf)}{2nf} \sim D_{2nf} \cdot 2^{(n+p)} \text{ or } 2$ + $\frac{1}{2f} [D_{2nf} \times D^{(s)}] = 2^{(\alpha,\beta)} = \overline{2}$	$*[e_{2nf}^{\alpha}, e_{2nf}^{\alpha}]$ or $*[1, j]$	See	
$\pm \frac{1}{f} [C_{nf} \times C^{(s)}] \cdot 2^{(\gamma)} $ or 2	* $[e_{2nf}, e_{2nf}]$ or * $[1, j]$	in	
$+\frac{1}{f}[C_{nf} \times C_{nf}^{(s)}] \cdot 2^{(\gamma)}$	*[1, c_{2nf}] *[1, $e^{\gamma(f,s+1)}$]	Appendix	
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