

Examples of Diophantine relations:

1.  $a \equiv b \pmod{c}$ :  $\exists x: a = cx + b$  or  $b = cx + a$ . (over  $\mathbb{N}$ !)
2.  $a \geq b$ :  $a = b + x$  ( $\exists$  is always implicit)
3.  $a > b$ : (Exercise)
4.  $a = b \pmod{c}$ :  $a \equiv b \pmod{c}$  and  $0 \leq a < c$ .
5.  $\{(a, b, c) \mid a = b^c\}$ ?
6.  $\{2^k\}$ ? Tarski believed not Diophantine

$$G_b(0) = 0, G_b(1) = 1, \quad \boxed{\begin{array}{l} G_n(n+1) = b \cdot G_b(n) - G_b(n-1) \\ G_b(n-1) + G_n(n+1) = b \cdot G_b(n) \end{array}} \quad \begin{array}{l} \text{close to the recursion for } b^{n-1} \\ \Leftrightarrow \text{symmetric between forward and backward} \end{array}$$

$$b = 4: [\dots, -15, -4, -1, ] 0, 1, 4, 15, 56, \dots$$

$$b = 3: 0, 1, 3, 8, 21, \dots$$

$$b = 2: 0, 1, 2, 3, 4, 5, \dots \text{ (useless?)}$$

**WE SHOW:**  $\boxed{a = G_b(c)}$  for (fixed)  $b \geq 3$  is Diophantine. ( $b \geq 4$  simplifies some arguments.)

Missing link! [ Yuri Matiyasevich 1970, Julia Robinson, Martin Davis, Hilary Putnam 1961 ]

The points  $\begin{pmatrix} x_n \\ y_n \end{pmatrix} = \begin{pmatrix} G_b(n+1) \\ G_b(n) \end{pmatrix}$  lie on the hyperbola  $\boxed{h_b(x, y) := x^2 - bxy + y^2 - 1 = 0}$ . ( $\rightarrow$  picture)

**Lemma 1.** *The only integer solutions of  $h_b(x, y) = 0$  with  $x > y \geq 0$  are those points.*

Proof: The hyperbola is invariant under the shift  $\begin{pmatrix} G_b(n) \\ G_b(n-1) \end{pmatrix} \leftrightarrow \begin{pmatrix} G_b(n+1) \\ G_b(n) \end{pmatrix}$ :  $\begin{pmatrix} x \\ y \end{pmatrix} \mapsto \begin{pmatrix} bx-y \\ x \end{pmatrix}$  or  $\begin{pmatrix} x \\ y \end{pmatrix} \mapsto \begin{pmatrix} y \\ by-x \end{pmatrix}$ , and the shift preserves  $y < x$ .  $\square$

Now we can generate  $\{G_b(n)\} = \{x \mid \exists y: h_b(x, y) = 0\}$  but we don't know  $n$ .

**Lemma 2.**  $b \equiv b' \pmod{u} \implies G_b(n) \equiv G_{b'}(n) \pmod{u}$  (Induction. Easy.)  $\square$

IDEA: Choose two appropriate moduli  $M$  and  $m$  to coordinate  $G_b(n)$  with  $n$ :

$$\begin{array}{ll} G_w(0), G_w(1), \dots, G_w(n), \dots \pmod{M} & w \equiv b \pmod{M} \quad \boxed{\rightarrow G_b(n)} \\ G_w(0), G_w(1), \dots, G_w(n), \dots \pmod{m} & w \equiv 2 \pmod{m} \quad \boxed{\rightarrow G_2(n) = n} \end{array}$$

As long as  $n$  is small and  $G_w(n) \leq M$ , we have  $G_w(n) \pmod{M} = G_b(n)$  (and  $G_w(n) \pmod{m} = n$ ), but for larger  $n$ ,  $G_w(n) \pmod{M}$  gets out of control, and there will be extra solutions. ( $\rightarrow$  picture)

- Make  $G_w(n) \pmod{M}$  mirror-symmetric after reaching a peak at  $G_w(p) = G_b(p)$ : ( $\rightarrow$  picture)  
 $G_b(p-1) \equiv G_b(p+1) \pmod{M} \implies M := G_b(p+1) - G_b(p-1)$  ( $p = \text{peak} = \text{period}$ )
- Avoid “negative” values by using the absmod operation instead of mod. ( $\rightarrow$  picture)  
 $x \text{ absmod } M = a \iff x = qM \pm a$  and  $0 \leq a \leq M/2$ .
- The period  $m$  of “ $G_w(n) \text{ absmod } m = n$ ” should divide the period  $2p$  of “ $G_w(n) \text{ absmod } M$ ”:  $m \mid p$ .
- Choose  $m$  (and  $M$ ) larger than twice the (supposed) value  $a$  of  $G_b(c)$ , so that absmod does no harm.

1.  $m > 2a$ .
2.  $p$  should be a multiple of  $m$
3.  $M := G_b(p+1) - G_b(p-1)$
4. Choose  $w > 2$  with  $w \equiv b \pmod{M}$   
 $w \equiv 2 \pmod{m}$
5.  $h_w(x, y) = 0$  [  $\implies x = G_w(n)$  for some  $n$  ]
6.  $a = x \text{ absmod } M$  [  $a = G_b(n)$  ]  
 $c = x \text{ absmod } m$  [  $c = G_2(n) = n$  ]

**Lemma 3.**  $G_b(k)^2 \mid G_b(p) \implies G_b(k) \mid p$

(Cf. Fibonacci numbers:  $F_k \mid F_p \iff k \mid p$ .)

Application: Choose  $m$  of the form  $m = G_b(k)$  for some  $k$ , by requiring  $\boxed{h_b(m, m') = 0}$ .

Then  $m^2 \mid G_b(p) \implies m \mid p$ .

Implementation of Conditions 2 and 3.

$$\boxed{h_b(r, s) = 0, r < s} \quad \begin{aligned} G_b(p-1) &= r, \text{ for some } p \\ G_b(p) &= s \\ G_b(p+1) &= bs - r \end{aligned}$$

$$\boxed{M = (bs - r) - r} \quad [ = G_b(p+1) - G_b(p-1). \text{ Also } G_b(p) < M/2. ]$$

$$\boxed{m^2 \mid s} \quad [ \implies m \mid p. ]$$

The conditions are enough to ensure that every solution  $(a, c)$  satisfies  $a = G_b(c)$ . (The condition  $a < m/2$  cuts off extra solutions.)

Converse direction:

We need to show that  $m, M$  with  $\gcd(m, M) = 1$  exist (then  $w$  satisfying (4.) exists, by the Chinese Remainder Theorem), and that  $m^2 \mid s$  can be fulfilled.

- Choose  $m = G_b(k)$  odd and set  $s = G_b(p)$  for  $p = k \cdot m$ :

Then it can be shown that  $\gcd(m, M) = 1$  and  $m^2 \mid s$ .

Getting to the relation  $a = b^c$ :

$$(b-1)^n \leq G_b(n+1) \leq b^n$$

$$b^c = \lim_{x \rightarrow \infty} \frac{G_{bx+4}(c+1)}{G_x(c+1)} \approx \frac{(bx \pm \text{const})^c}{(x \pm \text{const})^c} \rightarrow b^c, \text{ for all } b, c \geq 0$$

$$b^c = \left\lfloor \frac{G_{bx+4}(c+1)}{G_x(c+1)} \right\rfloor \text{ for } x > 16(c+1)G_{b+4}(c+1).$$

(The “+4” term ensures that this works even for  $b = 0$ .)