# Discrete Geometry, SS 2011 - exercise sheet 8+9 

due date: Tuesday, June 14th, 2011, 14:00
due date for exercises 42-45: Tuesday, June 21st, 2011, 14:00
35. Tightness of the shatter function lemma (10 credits).

Show that the bound

$$
\Phi_{d}(m)=\binom{m}{0}+\binom{m}{1}+\binom{m}{2}+\cdots+\binom{m}{d}
$$

on the shatter function is tight, by constructing a range space of VC-dimension $d$ with a ground set of $m$ elements and $\Phi_{d}(m)$ ranges.
(Hint: An easy solution can be guessed by staring at the formula. Can you find different solutions, for example for $d=1$ ? For $d=2$ ?)
36. Preserving the VC-dimension (10 credits)

Show that the operation of taking the symmetric difference $\oplus$ with a fixed set $A \subseteq X$ does not change the VC-dimension of a range space $(X, \mathcal{F})$ :

$$
\mathcal{F}^{\prime}:=\{R \oplus A \mid R \in \mathcal{F}\}
$$

Can it change the shatter function? Can it change the required size of $\varepsilon$-nets? Can it change the required size of $\varepsilon$-approximations?
37. ( 0 credits) Show that in a range space with VC-dimension $d=1$, there is always an $\varepsilon$-net with at $\operatorname{most} \max \left\{2,\left\lceil\frac{1}{\varepsilon}\right\rceil-1\right\}$ elements, and this bound is tight.
38. ( 0 credits) Prove that $\Phi_{d}(m) \leq(e m / d)^{d}$ for all $d \geq 1$ and $m \geq 1$.
39. Iterated $\varepsilon$-approximations ( 10 credits)
(a) Prove: If $N_{1}$ is an $\varepsilon_{1}$-approximation of a finite range space $(X, \mathcal{F})$ and if $N_{2}$ is an $\varepsilon_{2}$-approximation of the induced range space $\left(N_{1},\left.\mathcal{F}\right|_{N_{1}}\right)$, then $N_{2}$ is an $\left(\varepsilon_{1}+\varepsilon_{2}\right)$ approximation of the original range space $(X, \mathcal{F})$.
(b) If $N_{1}$ is only an $\varepsilon_{1}$-net (and not an $\varepsilon_{1}$-approximation), can one still conclude that $N_{2}$ is at least an $\left(\varepsilon_{1}+\varepsilon_{2}\right)$-net?
(c) If, on the other hand $N_{1}$ is an $\varepsilon_{1}$-approximation but $N_{2}$ is only an $\varepsilon_{2}$-net, can one conclude that $N_{2}$ is an $\left(\varepsilon_{1}+\varepsilon_{2}\right)$-net?
40. Deviation from the mean ( 0 credits)

Let $X=X_{1}+X_{2}+\cdots+X_{n}$ be the sum of $n$ independent Bernoulli random variables with success probability $P\left[X_{i}=1\right]=p$ (and $\left.P\left[X_{i}=0\right]=1-p\right)$. Show that $P[X \geq$ $\left.\frac{1}{2} n p\right] \geq \frac{1}{2}$, provided that $n p \geq 8$, by applying the Chebysheff inequality, which says that $P[|X-\mu| \geq t \sigma] \leq 1 / t^{2}$, for any random variable $X$ with mean $\mu$ and variance $\sigma$.
41. VC-dimension of half-spaces (0 credits)

Show that the range space of all half-spaces $\left\{a_{1} x_{1}+a_{2} x_{2}+\cdots+a_{d} x_{d} \leq b\right\}$ of $\mathbb{R}^{d}$ has VC-dimension $d+1$.

Hint for the upper bound: Take any $d+2$ points $p_{1}, \ldots, p_{d+2}$ in $\mathbb{R}^{d}$. By rank considerations, the system of equations $\sum \lambda_{i} p_{i}=0, \sum \lambda_{i}=0$, has a nonzero solution $\left(\lambda_{1}, \ldots, \lambda_{d+2}\right)$. Then the subset $\left\{p_{i} \mid \lambda_{i}>0\right\}$ cannot be cut out by a half-space.
42. Lifting (10 credits)
(a) Give a finite upper bound on the VC-dimension of the ranges of $\mathbb{R}^{2}$ that can be defined by quadratic equations $a x^{2}+b x y+c y^{2}+d x+e y+f \leq 0$, by using the mapping $(x, y) \mapsto\left(x^{2}, x y, y^{2}, x, y\right)$ from $\mathbb{R}^{2}$ to $\mathbb{R}^{5}$ (the so-called Veronese map) and applying the previous exercise.
(b) Extend the result to $d>2$ dimensions.
(c) Show that, for balls in $\mathbb{R}^{d}$, the general bound can be improved by using the lifting $\operatorname{map}\left(x_{1}, \ldots, x_{d}\right) \mapsto\left(x_{1}, \ldots, x_{d}, x_{1}^{1}+\cdots+x_{d}^{2}\right)$.
(d) (0 credits) Show that, for balls in $d$ dimensions, the shatter function has actually the stronger bound $\Phi_{d+1}(n)$, and this is tight for $n$ points in general position.
43. ( 0 credits) Show that the unit square cannot be expressed as $\{(x, y) \mid p(x, y) \geq 0\}$ with a single polynomial $p(x . y)$.
44. Finite VC-dimension under Boolean operations (8 credits)

Suppose that $(X, \mathcal{F})$ is a range space of VC-dimension 5 , and $(X, \mathcal{G})$ is a range space of VC-dimension 9. Find a finite bound on the VC-dimension of the range space

$$
(X,\{R \cap S \mid R \in \mathcal{F}, S \in \mathcal{G}\})
$$

Hint: try to bound the shatter function.
45. $\varepsilon$-nets and the cutting lemma (12 credits)
(a) Let $H$ be a finite set of lines in the plane. For a triangle $T$, let $H_{T}$ be the set of lines intersecting the interior of $T$. let $\mathcal{T} \subseteq 2^{H}$ be the system of sets $H_{T}$ for all triangles $T$. Show that the VC-dimension of $\mathcal{T}$ is bounded by a constant.
(b) With the $\varepsilon$-net theorem, prove a weak form of the cutting theorem: For every finite set $H$ of lines and every $r, 1<r<|H|$, there is a $\frac{1}{r}$-cutting for $H$ consisting of $O\left(r^{2} \log ^{2} r\right)$ triangles (or trapezoids).
46. Vertex neighborhoods (0 credits)

For a given undirected graph $G=(V, E)$, let $N(v)=\{u \in V \mid u v \in E\}$ denote the neighborhood of the vertex $v \in V$ (not including $v$ itself). Show that there is a constant $c$ such the set system $\{N(v) \mid v \in V\}$ of vertex neighborhoods of any planar graph has VC-dimension bounded by $c$.

