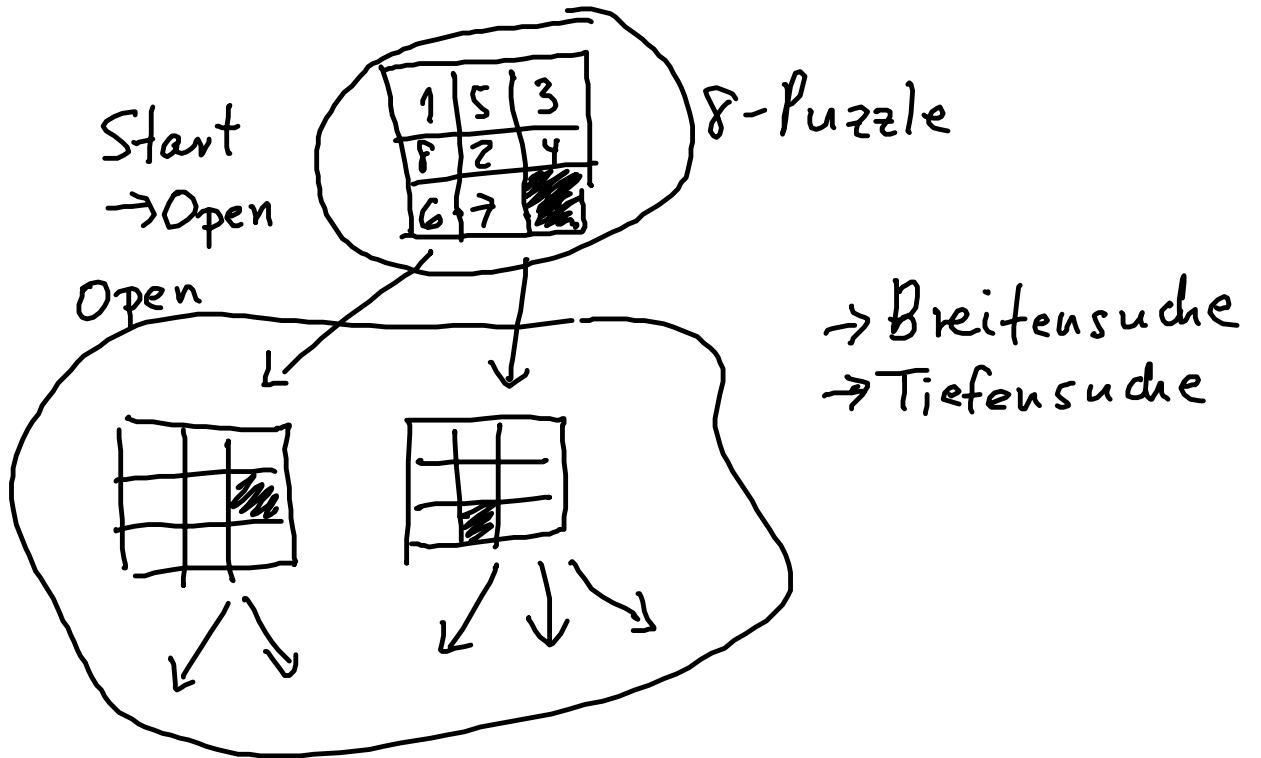
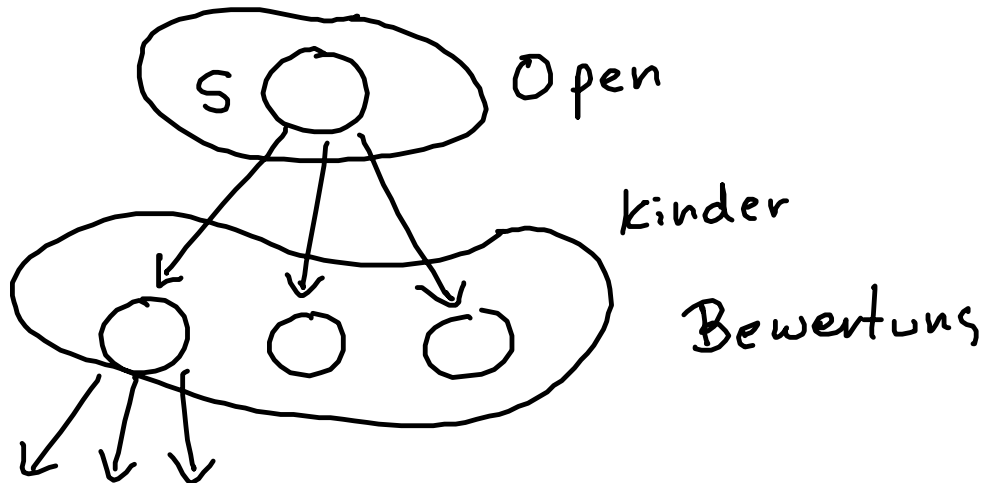


Zustandsraum



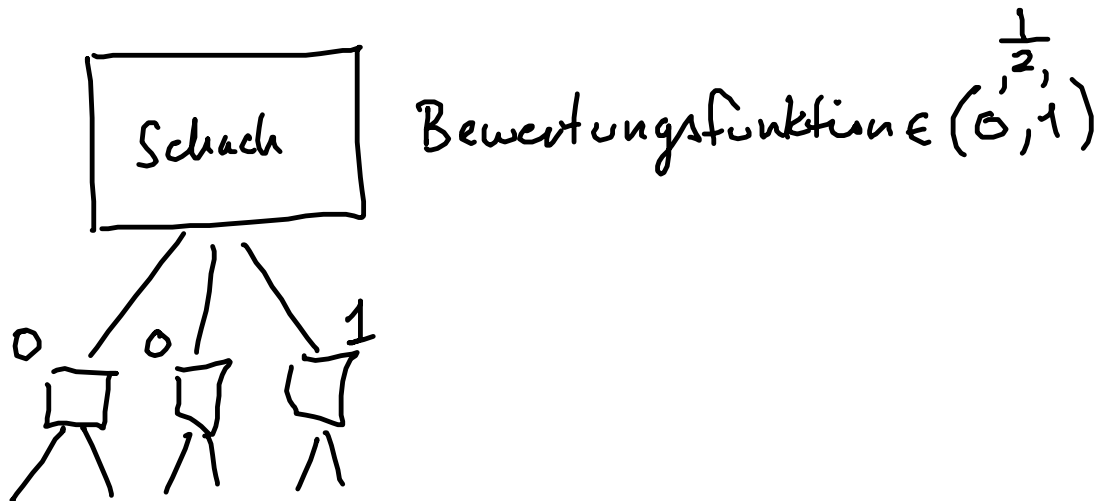
Bestensuche



$$\text{New Open} := (\text{Open} - S) \cup \text{Kinder}$$

sortiert
 "Maß
 Abstandsmaß"

→ Heuristik



Bewertung
 1 + 1 + 0
 + 1 + 0 + 1
 + 1 + 1 + 1
7 Fehler

Heuristik-1

2 Fehler

2	1	3
6	5	7
8	 	4

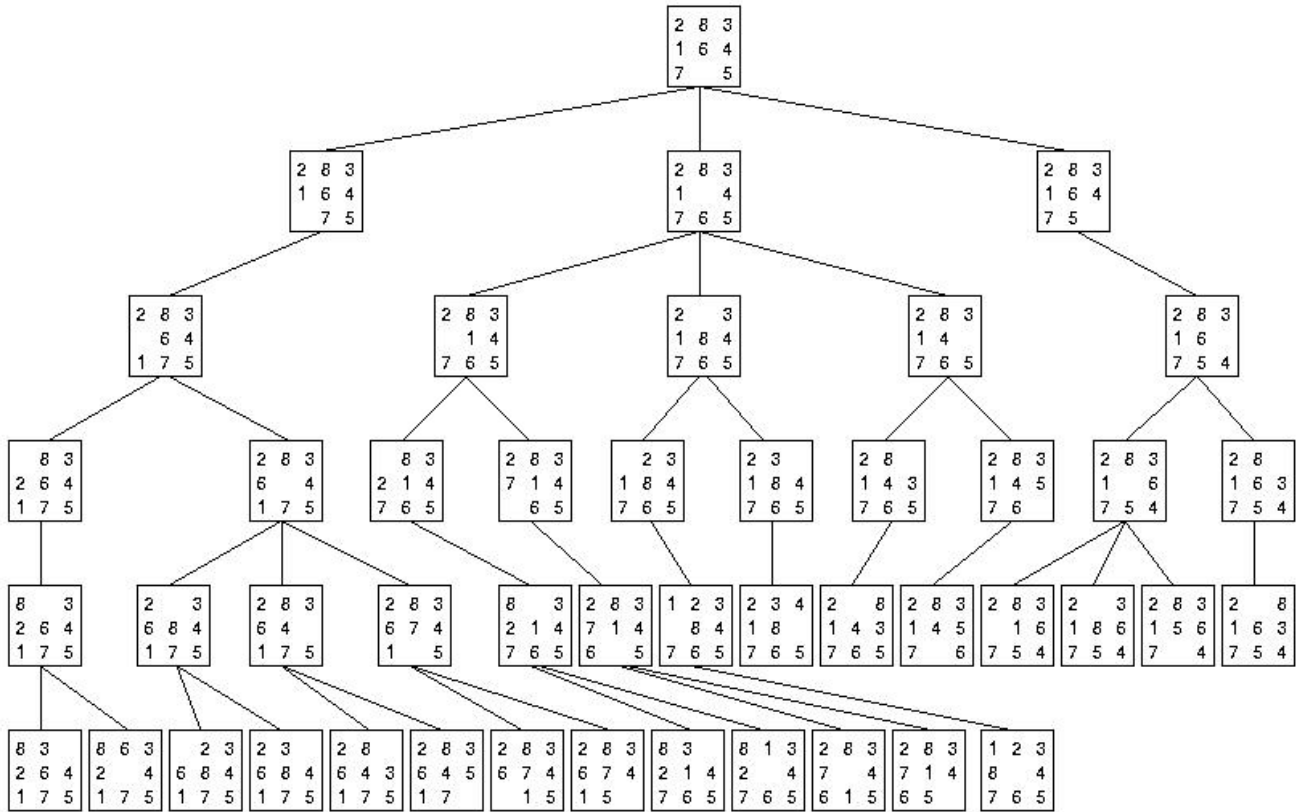
1	2	3
4	5	6
7	8	

2	1	3
4	5	6
7	8	

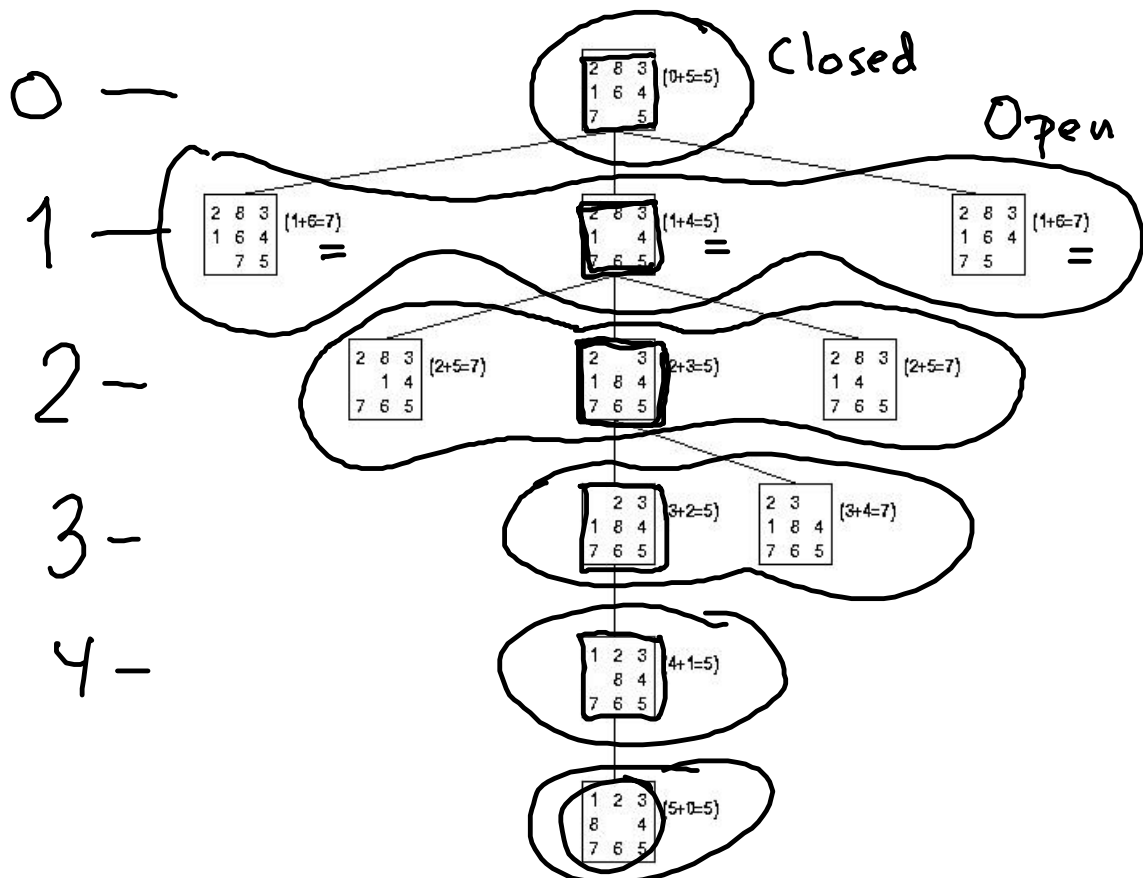
Heuristik-2

1 + 1 + 0
 + 2 + 0 + 3
 + 1 + 3 + 1
12

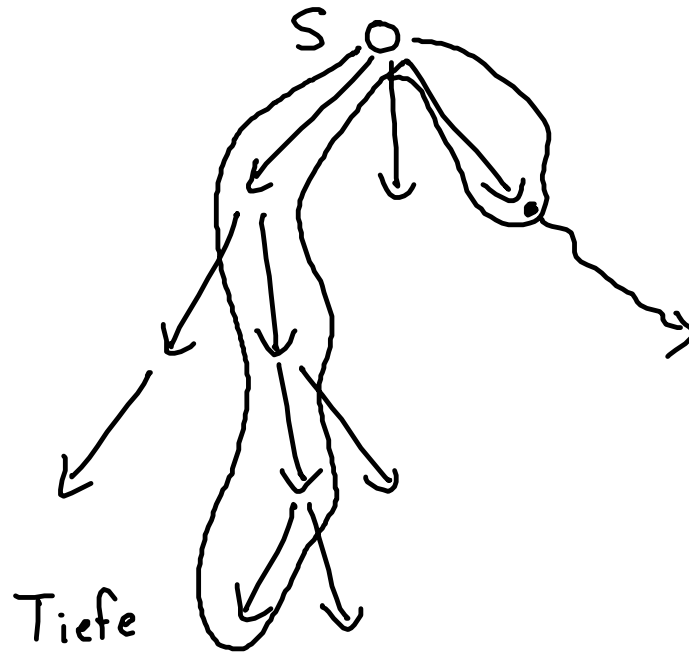
Bestensuche



Heuristik

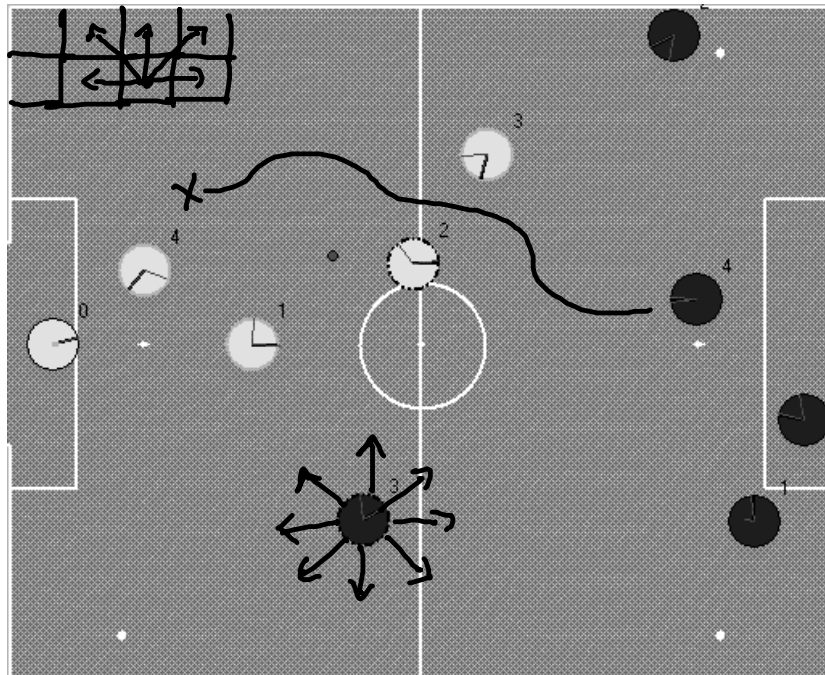


$$f(n) = g(n) + h(n)$$



$$f(n) = g(n) + h(n)$$

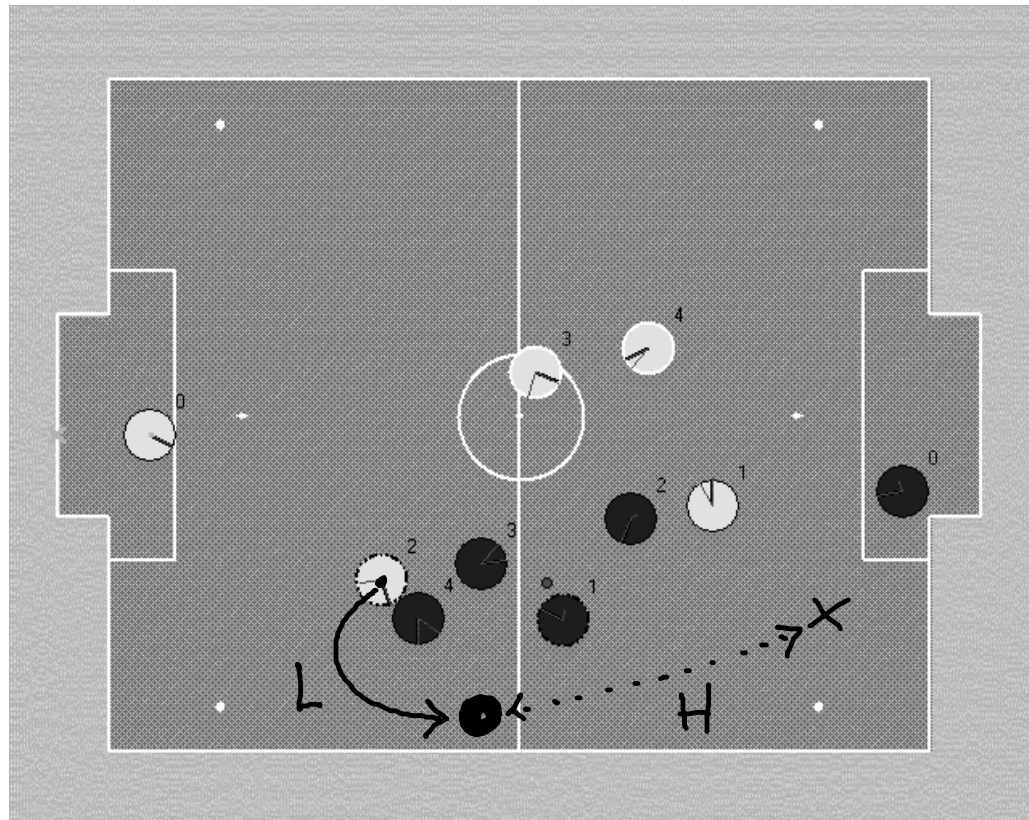
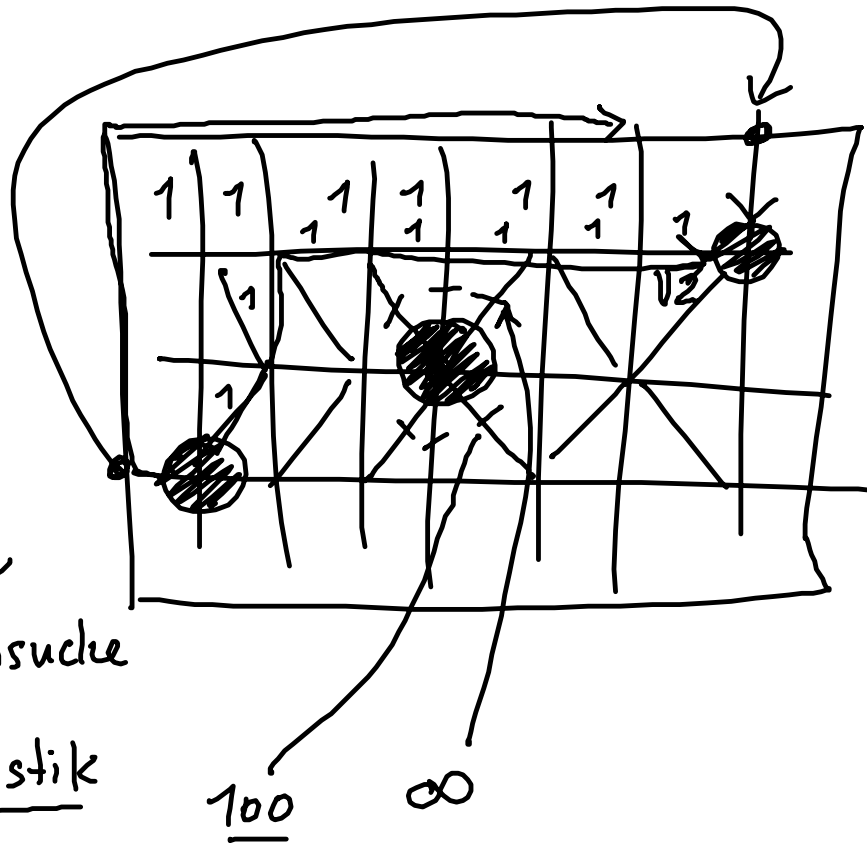
zu groß



Dijkstra

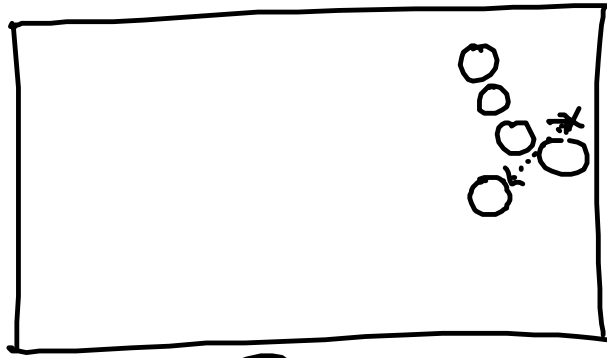
A*

Bestensuche
+
Heuristik

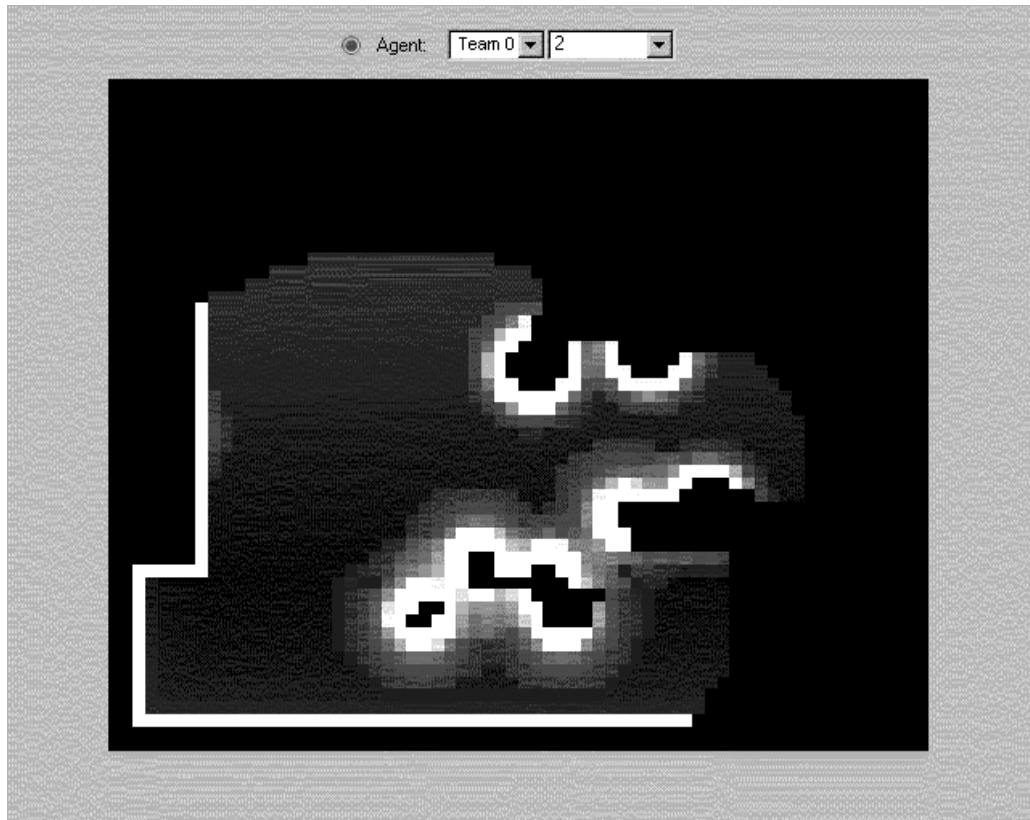
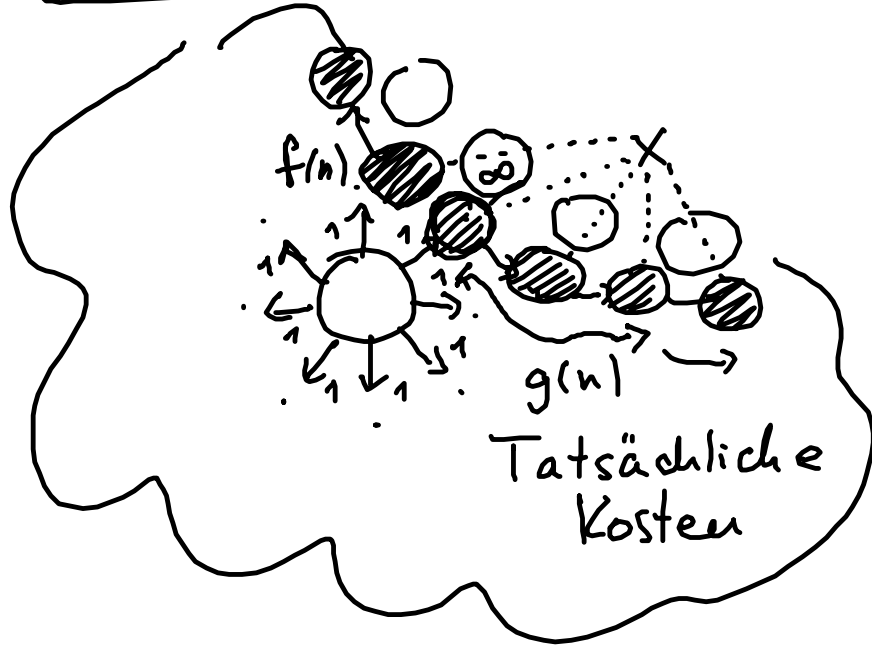


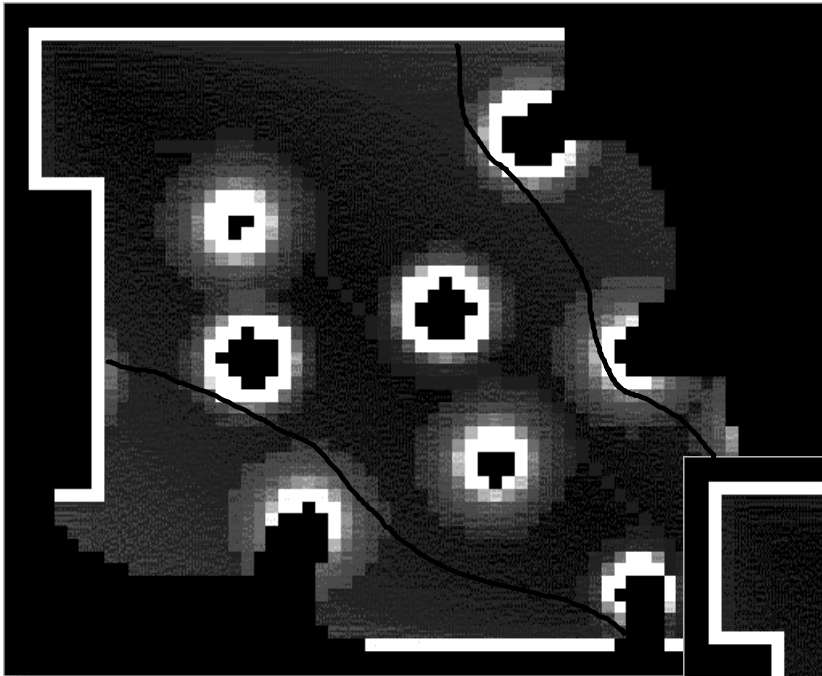
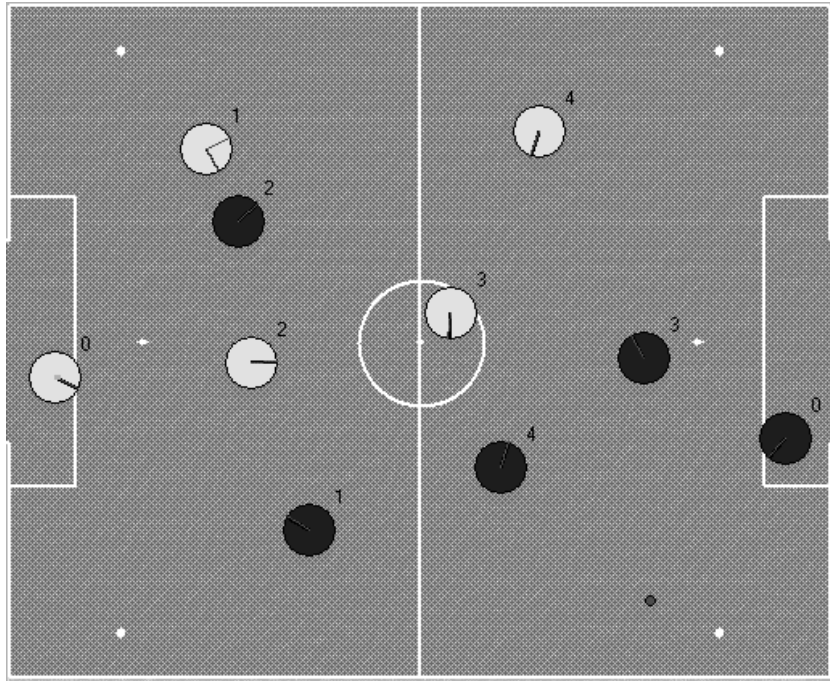
$$f(n) = g(n) + h(n)$$

L H

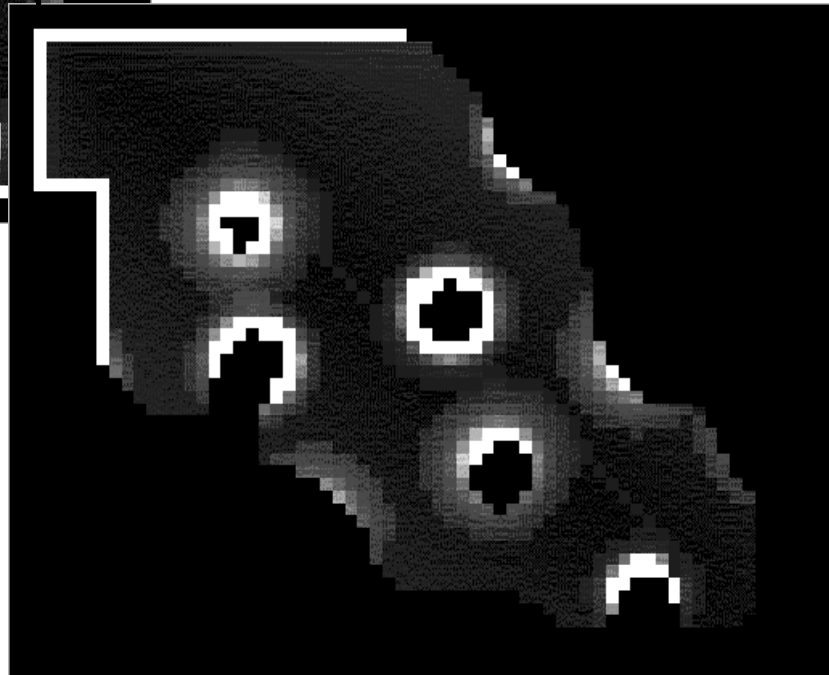


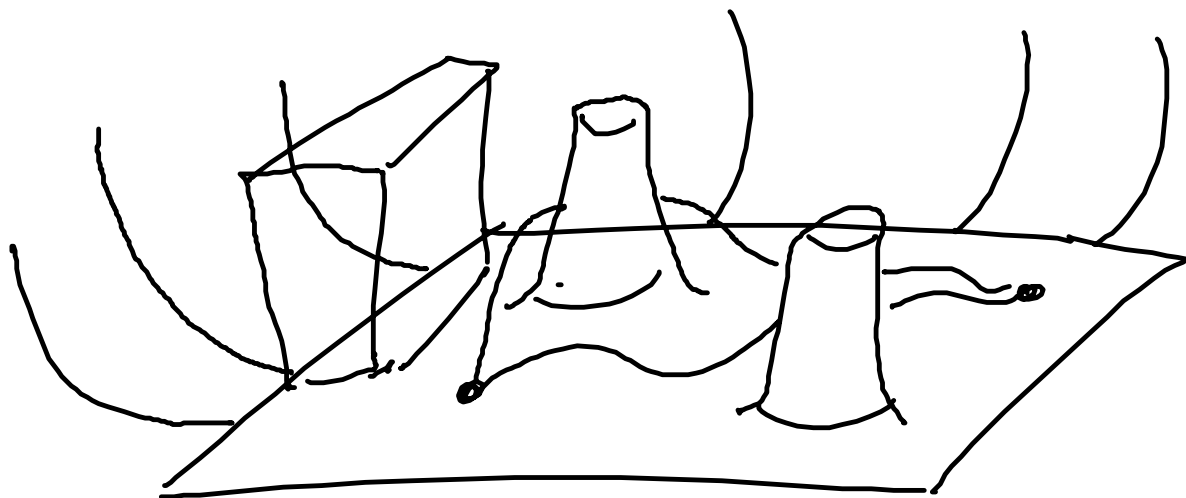
H





Potential
Methode





Sortieren

Repräsentation

$bd(2, 3, 1, 5, 6, 7, 8, 4, 6)$

Bewertung

$$f(n) = g(n) + h(n)$$

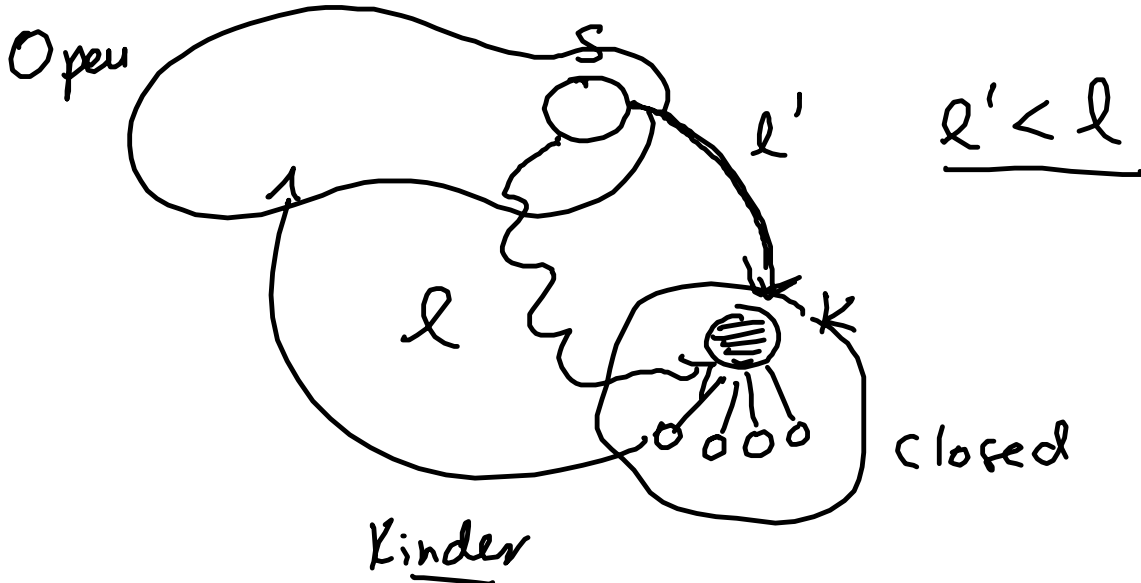
~
Tiefe

Anzahl der
falsch platzierte
Steine

Liste von Paaren

Open $[[bd(\dots), 5], [bd(\dots), 6] \dots]$
 Butt $f(n)$

Closed $[[bd(\dots), 7], [bd(\dots), 8] \dots]$



Goals Open Closed

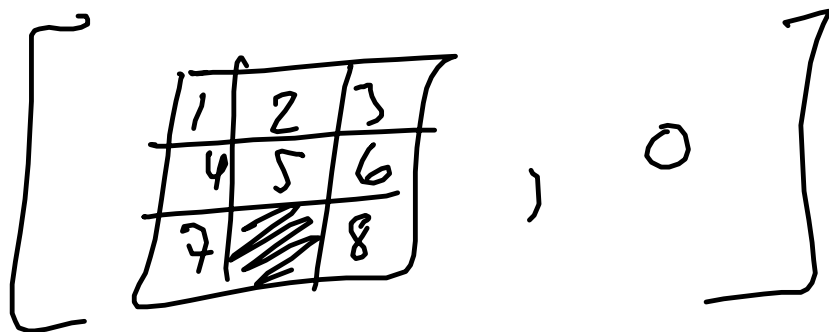
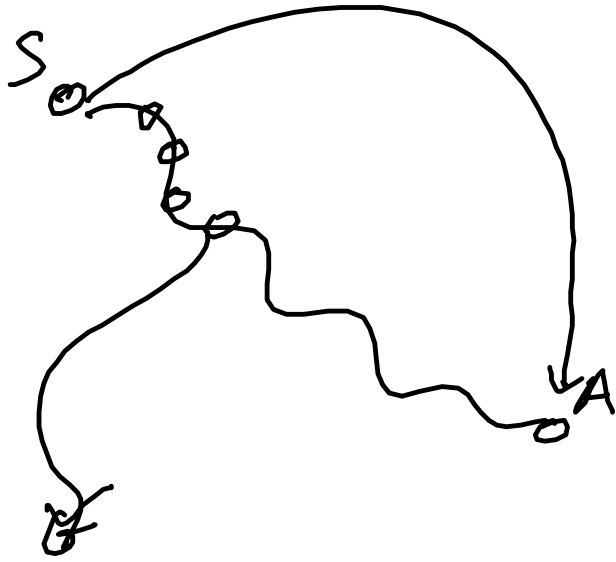
best (Goals, [[X, -] | -], -) :-
 member(X, Goals), !, write(X).

best (Goals, [[X, T] | RestOpen], Closed) :-
 member([X, T1], Closed),
 T >= T1, !,
 best (Goals, RestOpen, Closed).

best (Goals, [[X, T] | RestOpen], Closed) :-
 member([X, T1], Closed),

$T < T1, !,$
kinderbest (X, T, L),
appendm (L, RestOpen, NewOpen),
qsort (NewOpen, Open),
best (Goals, Open, [[X|T]|
Closed]).

best (Goals, [[X|T]|RestOpen], Closed):-
write('open '), write(X),
heuristik (X, T, FX), write(FX),
kinderbest (X, T, L),
appendm (L, RestOpen, NewOpen),
qsort (NewOpen, Open),
best (Goals, Open, [[X|T]|Closed]).



mygsort () ← selber programmieren

kinderbest(x, T, L) :-

T1 is T+1,

setof([Y, T1], move(x, Y), L), !.

kinderbest(-, -, []).

_____ 0 _____

Bestensuche

A^*

$$f(n) = g(n) + h^*(n)$$

↑ ↑
Tiefe Heuristik

Echte Kosten $h(n)$

$$\underbrace{h^*(n) \leq h(n)}$$

Heuristik die Aufwand unterschätzt

A^* ist Bestensuche mit einer Heuristik, die den echten Aufwand unterschätzt.

⇒ findet eine Lösung, falls sie existiert (admissible)

Fall

1

~~2~~

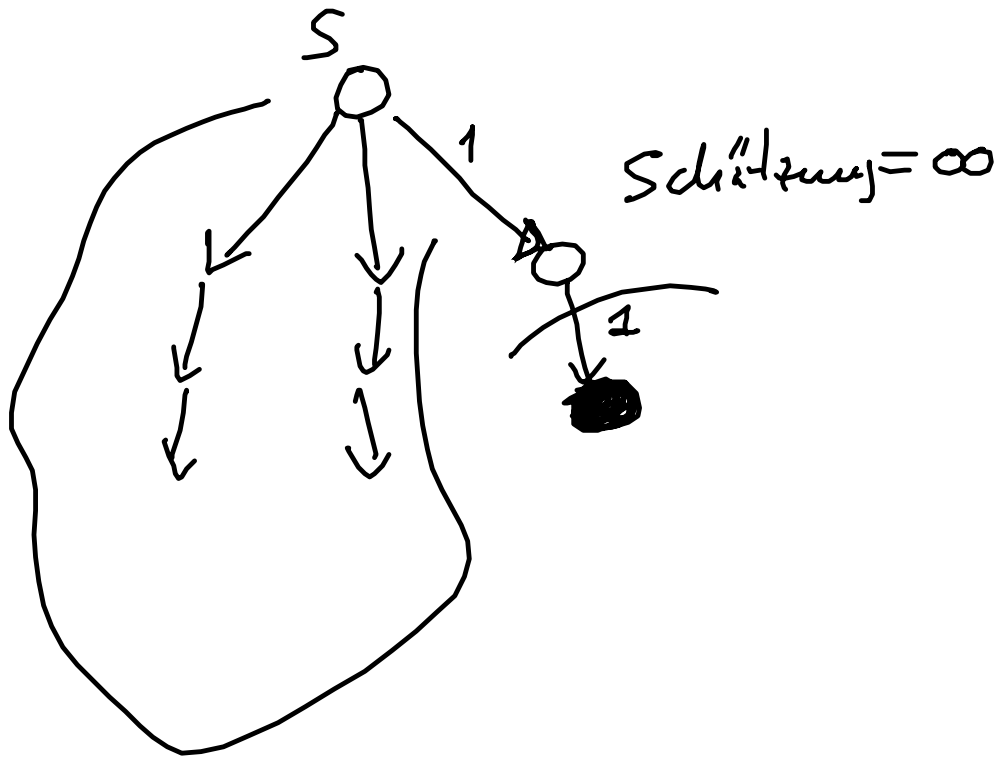
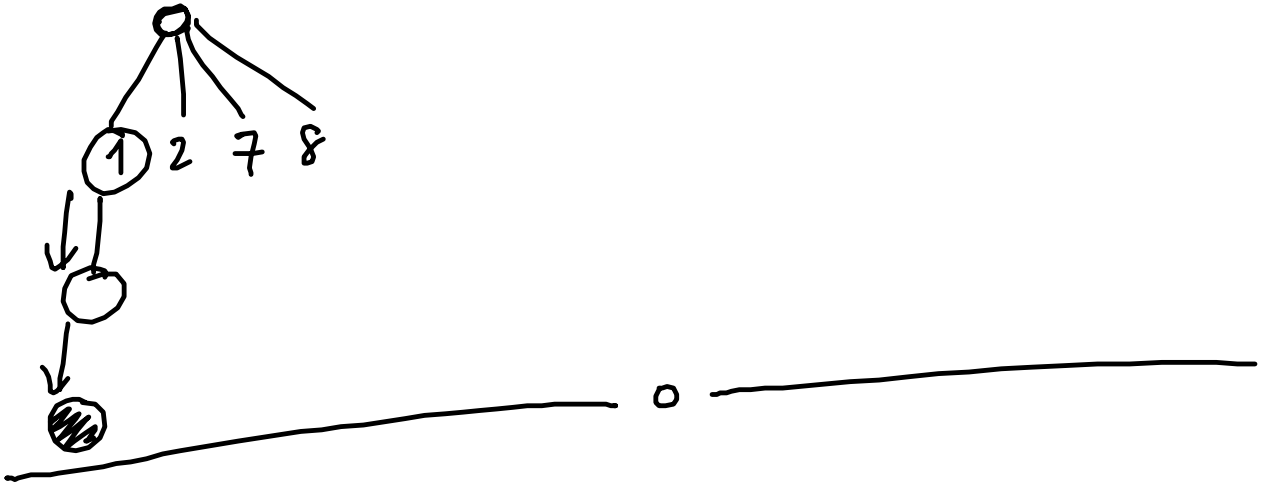
2

3

$$h^*(n) = h(n)$$

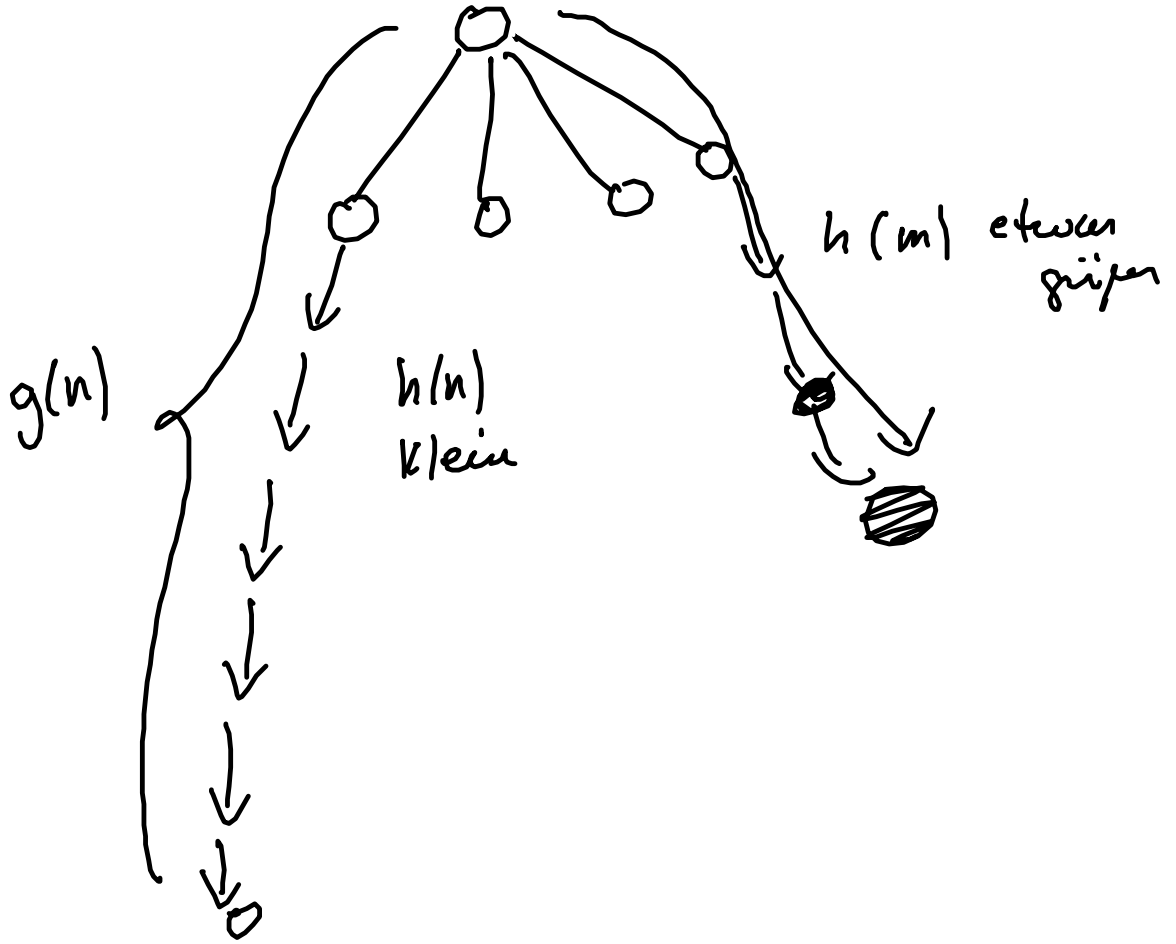
$$h^*(n) \leq h(n)$$

$$h^*(n) > h(n)$$



Unterschätzung

S



letztes Bsp.

