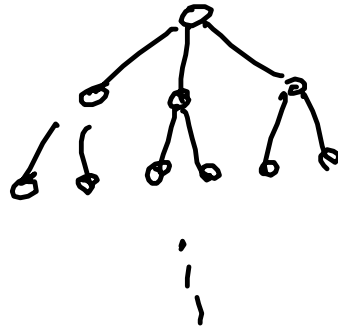


# Suchraumproblematik

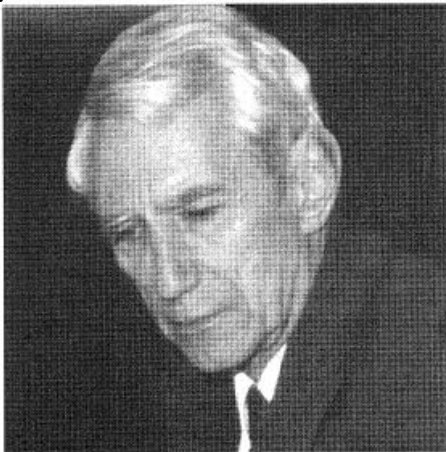
Schach



Verzweigungsfaktor ca. 35-40

Halbzugtiefe	#Stellungen
1	40
2	1600
⋮	
7	164 Milliarden
⋮	
11	419 Billionen ↑

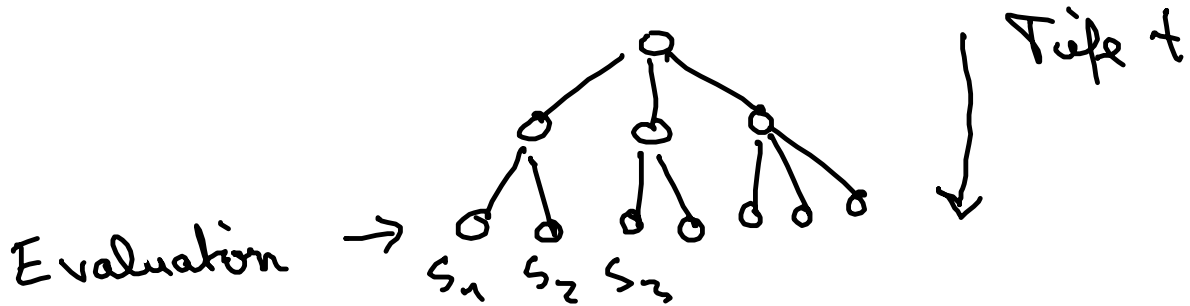
## Claude Elwood Shannon



- Mathematiker, New Jersey
- Vortrag, 9. März 49  
→ Min Max + Evaluation
- Zurechnungstheorie

# IDEE

- feste Suchtiefe
- Abschätzung

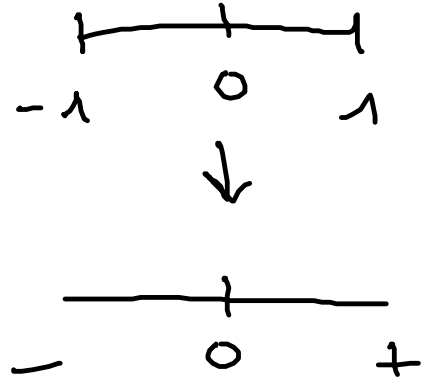


## Bewertungsfunktion

meistens linear

$S$  zu beurteilende Stellung

$$f: S \rightarrow \mathbb{R}$$

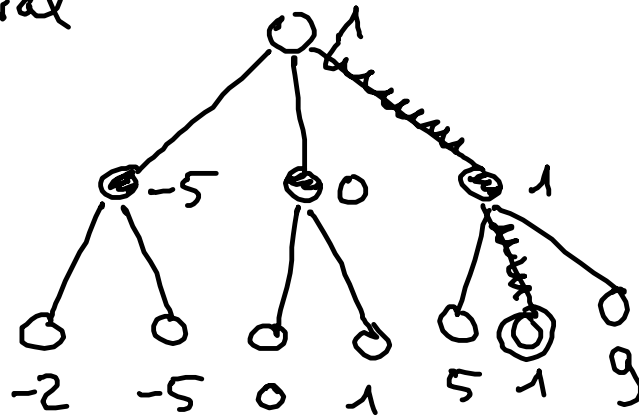


$$f(s) = \sum_{i=1}^n \alpha_i \cdot f_i(s)$$

Spiel-würsen

$$\left\{ \begin{array}{l} f_1(s) = \text{material}(W_s) - \text{material}(S_s) \\ f_2(s) = \text{mobilität}(W_s) - \text{material}(S_s) \\ \vdots \end{array} \right.$$

Min Max + Eval



PseudoCode

max Knoten (Knoten x)  $\rightarrow$  Integer  
 if (x ist Blatt)

```

return evaluation(x)
else { //  $x_1 \dots x_k$  seien Kinder von  $x$ 
 $w := -\infty$ 
for  $i := 1$  to  $k$  do {
 $v := \text{minKnoten}(x_i)$ 
if ( $v > w$ ) then  $w := v$ 
}
return  $w$ 
}

```

$\text{minKnoten}(\text{Knoten } x) \rightarrow \text{ZuInteger}$   
 if ( $x$  Blatt)

```

return evaluation(x)
else {
 $w := +\infty$ 
for  $i := 1$  to  $k$  do {
 $v := \text{maxKnoten}(x_i)$ 
if ( $v < w$ ) then  $w := v$ 
}
return  $w$ 
}

```

# Min Max in Prolog

minmax (Pos, BestSucc, Val) :-  
 moves (Pos, PosList), !,  
 best (PosList, BestSucc, Val);  
 evaluation (Pos, Val).

best ([Pos], Pos, Val) :-  
 minmax (Pos, \_, Val).

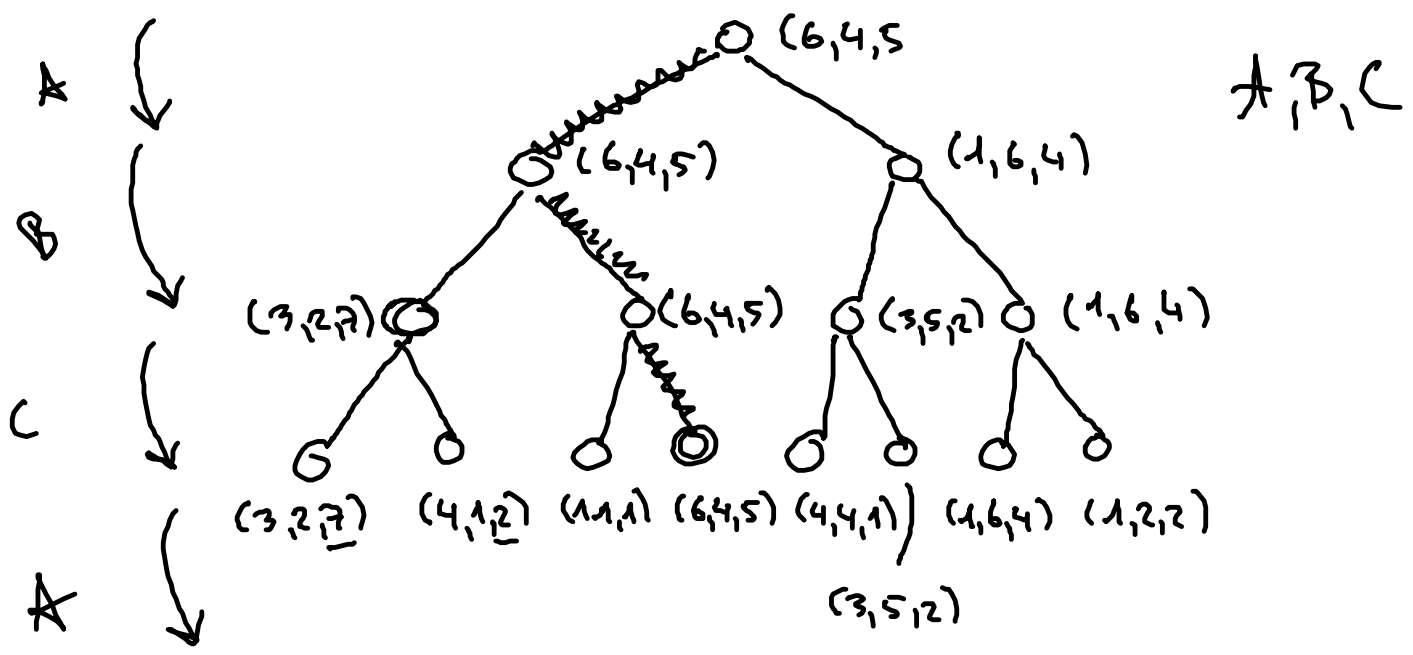
best ([Pos1 | PosList], BestPos, BestVal) :-  
 minmax (Pos1, \_, Val1),  
 best (PosList, Pos2, Val2),  
 betterof (Pos1, Val1, Pos2, Val2,  
 BestPos, BestVal).

betterof (Pos0, Val0, Pos1, Val1, Pos0, Val0) :-  
 min\_to\_move (Pos0),  
 Val0 > Val1, !,  
 ;  
 max\_to\_move (Pos0),  
 Val0 < Val1, !.

betterof (Pos0, Val0, Pos1, Val1, Pos1, Val1).

Ausblick

Minimax bei Mehr Spieler Spielen



(X, 4, 2)  
 ↑ ↑ ↑  
 A B C

z.B. Siedler

→ Allianzen (Zusammenschlüsse)

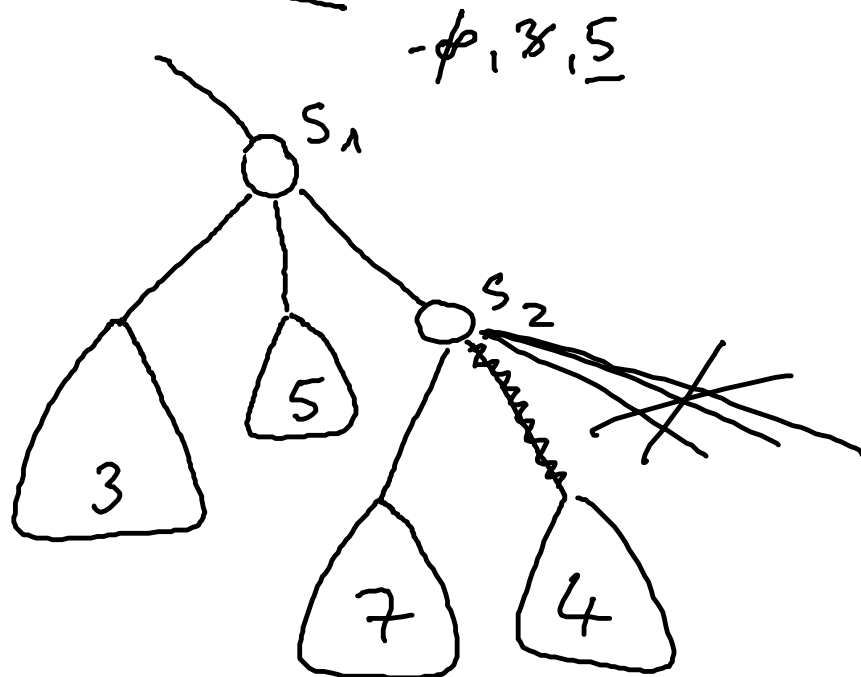
Alan Newell, John Shaw und  
Herbert Simon

- „Computerwissenschaftler“, Pittsburgh

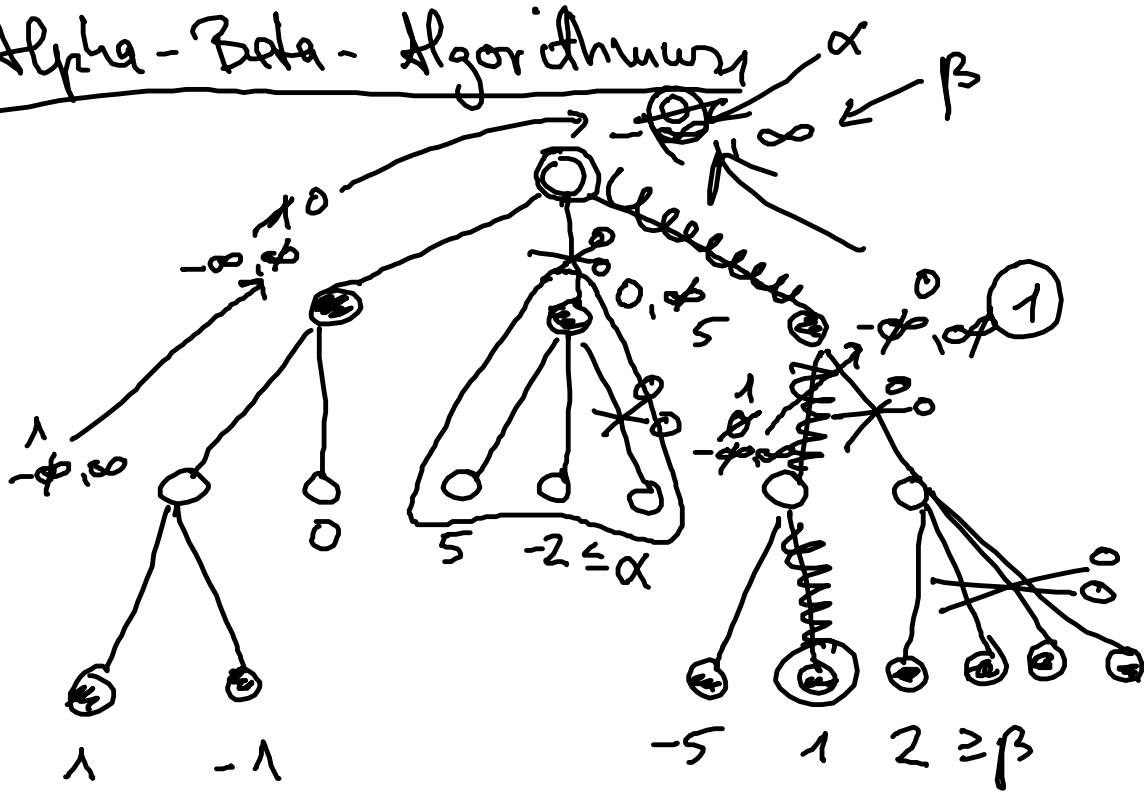
- 1958

→ Alpha-Beta

lokale Betrachtung



# Alpha-Beta-Algorithmus



## Pseudocode

$\alpha\beta$ max Knoten(Knoten  $x$ , Integer  $\alpha, \beta$ )  $\rightarrow$  Integer  
 if ( $x$  ist Blatt)

    return evaluation( $x$ )

else {

$w := \alpha$

    for  $i := 1$  to  $k$  do {

$v := \alpha\beta$ min Knoten( $x_i, w, \beta$ )

        if ( $v > w$ ) then  $w := v$

        if ( $w \geq \beta$ ) then return  $w$  //  $\beta$ -Cuttoff

    }

    return  $w$



```

}
αβ min Knoten (Knoten x, Integer α, β) → Integer
if (x ist Blatt)
    return ebaluator(x)
else {
    ω := β
    for i := 1 to k do {
        v := αβ max Knoten(xi, α, ω)
        if (v < ω) then ω := v
        if (ω ∈ α) then return ω
            // α-Cutoff
    }
    return ω
}
}

```

## Effektivität von Alpha-Beta

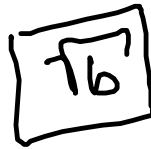
im schlechtesten Fall wie Min Max

### Zugreihenfolge

optimal

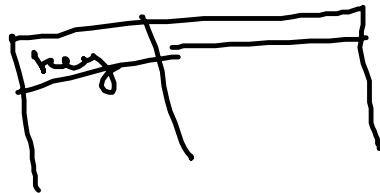
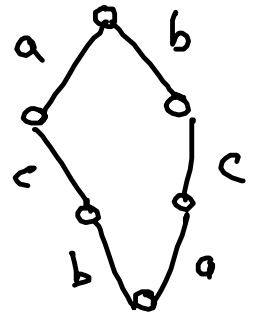
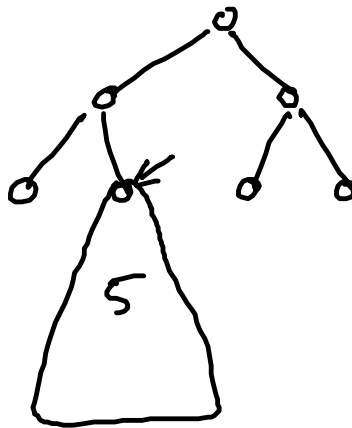
Bsp Schach

Verzweigungsgrad  
(branch factor b)



## Optimierung

### o Transpositionstabellen

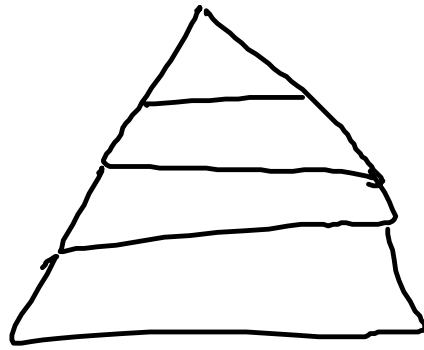


### o ZugSortierung

Schritt: Killer moves

History move

Hauptvariante



↓ 2, 3, 4

iteratives Suchverfahren