

# Hendricus van der Holst

Supervisor: Prof. dr. Martin Aigner  
Field of Research: Graph theory and Combinatorics  
Topic: graph polynomials and topological graph theory  
Postdoc Student: at the program since November 1 2000

## Field of Research and results

This semester we continued working on the graph polynomial introduced by Arratia, Bollobás, and Sorkin [1]; see last semester report. We discovered some new results, and a paper on this topic has been submitted.

One of the problems I am working on at the moment is to get a polynomial time algorithm for calculating the interlace polynomial for isotropic systems of bounded branch-width. Branch-width has been defined for graphs and matroids; for graphs it can be shown that the branch-width of the graph is small if and only if the tree-width of the graphs is small. It is straightforward to define branch-width for isotropic systems. A branch-decomposition of an isotropic system with  $n$  element consists of a ternary tree  $T$  (a tree in which each vertex has degree three or one) with exactly  $n$  leaves and a labelling of these leaves with the element of the isotropic system. For each edge  $e$  in  $T$ ,  $T \setminus e$  has exactly two components, and hence for each edge  $e$  of  $T$ , the set of elements of the isotropic system is partitioned into two sets  $X$  and  $Y$ , the elements corresponding to the leaves of one component and the elements corresponding to the leaves of the other component. Every partition  $(X, Y)$  of the set of elements of the isotropic system defines a separation of a certain order. So each edge of  $T$  defines a separation. The width of the branch-decomposition is defined as the maximum order over all these separations. The branch-width is defined as the minimum width over all branch-decompositions. It might also be interesting to characterize the isotropic systems that have branch-width bounded by small constant, say two or three.

I finally succeeded in making a new polynomial time algorithm for testing whether a graph has a linkless embedding or not. A linkless embedding of a graph  $G$  is an embedding of  $G$  into 3-space such that the images of every two disjoint circuits have zero linking number (taken over  $\mathbb{Z}_2$ ). The first polynomial time algorithm was given by Robertson and Seymour, and is a

consequence of their Graph Minors project. This is however a purely existence result. Later Robertson, Seymour, and Thomas [2] gave the excluded minors for the class of graphs that have a linkless embedding, making the algorithm explicit. However, their algorithm is more a theoretical result than that it can be used in practice. My algorithm seems to be more practical.

## Activities

I am currently at Oxford University. On 27 June I gave a talk about the interlace polynomial at the Combinatorial Theory Seminar, Oxford.

## Literatur

- [1] R. Arratia, B. Bollobás, and G.B. Sorkin, *The Interlace Polynomial: A New Graph Polynomial*, Proceedings of the Eleventh Annual ACM-SIAM Symposium on Discrete Algorithm, 237–245, San Francisco, California, 2000.
- [2] N. Robertson, P. Seymour, and R. Thomas, *Sach's Linkless Embedding Conjecture*, Journal of Comb. Theory, Series B, **64** (1995), 185–227.